

# Peierls bounds from Toom contours

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## Abstract

For deterministic monotone cellular automata on the  $d$ -dimensional integer lattice, Toom has given necessary and sufficient conditions for the all-one fixed point to be stable against small random perturbations. The proof of sufficiency is based on an intricate Peierls argument. We present a simplified version of this Peierls argument. Our main motivation is the open problem of determining stability of monotone cellular automata with intrinsic randomness, in which for the unperturbed evolution the local update rules at different space-time points are chosen in an i.i.d. fashion according to some fixed law. We apply Toom's Peierls argument to prove stability of a class of cellular automata with intrinsic randomness and also derive lower bounds on the critical parameter for some deterministic cellular automata.

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# Part I

## Results

### 1 Introduction and main results

#### 1.1 Introduction

Let  $\{0, 1\}^{\mathbb{Z}^d}$  denote the set of configurations  $x = (x(i))_{i \in \mathbb{Z}^d}$  of zeros and ones on the  $d$ -dimensional integer lattice  $\mathbb{Z}^d$ . By definition, a map  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  is *local* if  $\varphi$  depends only on finitely many coordinates, i.e., there exists a finite set  $\Delta \subset \mathbb{Z}^d$  and a function  $\varphi' : \{0, 1\}^\Delta \rightarrow \{0, 1\}$  such that  $\varphi((x(i))_{i \in \mathbb{Z}^d}) = \varphi'((x(i))_{i \in \Delta})$  for each  $x \in \{0, 1\}^{\mathbb{Z}^d}$ . We let  $\Delta(\varphi)$  denote the smallest set with this property, which may be empty: in this case  $\varphi$  is either constantly zero or one. We denote the constant functions by

$$\varphi^0(x) := 0 \quad \text{and} \quad \varphi^1(x) := 1 \quad (x \in \{0, 1\}^{\mathbb{Z}^d}). \quad (1.1)$$

A local map  $\varphi$  is *monotone* if  $x \leq y$  (coordinatewise) implies  $\varphi(x) \leq \varphi(y)$ . Let  $\{\varphi_0, \dots, \varphi_m\}$  be a set of monotone local maps  $\varphi_k : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$ , of which  $\varphi_0 = \varphi^0$  is the map that is constantly zero and  $\varphi_1, \dots, \varphi_m$  are not constant. Let  $\mathbf{r} = (\mathbf{r}(1), \dots, \mathbf{r}(m))$  be a probability distribution on  $\{1, \dots, m\}$ . We will be interested in i.i.d. collections of random variables  $\Phi^{p, \mathbf{r}} = \Phi^p = (\Phi_{i,t}^p)_{(i,t) \in \mathbb{Z}^{d+1}}$  with values in  $\{\varphi_0, \dots, \varphi_m\}$  such that

$$\mathbb{P}[\Phi_{i,t}^p = \varphi_k] = \begin{cases} p & \text{if } k = 0, \\ (1-p)\mathbf{r}(k) & \text{if } 1 \leq k \leq m, \end{cases} \quad (1.2)$$

where  $p \in [0, 1]$  is a parameter. We call  $\Phi^p$  a *monotone cellular automaton*. We will be interested in the case that  $p$  is small but positive. We think of  $\Phi^p$  as a small perturbation of  $\Phi^0$ . In the special case that  $m = 1$ , we say that  $\Phi^0$  is a *deterministic* monotone cellular automaton. If  $m \geq 2$  and  $\mathbf{r}(k) < 1$  for all  $k$ , then we say that  $\Phi^0$  has *intrinsic randomness*.

If  $X_0^p$  is a random variable with values in  $\{0, 1\}^{\mathbb{Z}^d}$ , independent of  $(\Phi_{i,t}^p)_{i \in \mathbb{Z}^d, t \in \mathbb{Z}_+}$ , then setting

$$X_t^p(i) := \Phi_{i,t}^p((X_{t-1}^p(i+j))_{j \in \mathbb{Z}^d}) \quad (i \in \mathbb{Z}^d, t > 0) \quad (1.3)$$

defines a Markov chain  $(X_t^p)_{t \geq 0}$  with state space  $\{0, 1\}^{\mathbb{Z}^d}$ . Let  $\mathbb{P}^x$  denote the law of this Markov chain started in a given initial state  $X_0^p = x$  and let  $\underline{0}$  and  $\underline{1}$  denote the configurations in  $\{0, 1\}^{\mathbb{Z}^d}$  that are constantly zero or one, respectively. It is well-known that

$$\mathbb{P}^{\underline{0}}[X_t^p \in \cdot] \xrightarrow[t \rightarrow \infty]{} \underline{\nu}_p \quad \text{and} \quad \mathbb{P}^{\underline{1}}[X_t^p \in \cdot] \xrightarrow[t \rightarrow \infty]{} \bar{\nu}_p \quad (1.4)$$

where  $\Rightarrow$  denotes weak convergence of probability measures on  $\{0, 1\}^{\mathbb{Z}^d}$ , equipped with the product topology, and  $\underline{\nu}_p$  and  $\bar{\nu}_p$  are invariant laws of the Markov chain defined in (1.3), that are called the *lower* and *upper* invariant laws, respectively. Let

$$\bar{\rho}_{\mathbf{r}}(p) = \bar{\rho}(p) := \lim_{t \rightarrow \infty} \mathbb{P}^{\underline{1}}[X_t^p(i) = 1] \quad (p \in [0, 1], i \in \mathbb{Z}^d) \quad (1.5)$$

denote the density of the upper invariant law, which by translation invariance does not depend on  $i \in \mathbb{Z}^d$ . Trivially,  $\bar{\rho}(0) = 1$ <sup>1</sup>. We say that the monotone cellular automaton  $\Phi^0$  defined by the monotone local maps  $\varphi_1, \dots, \varphi_m$  and the probability distribution  $\mathbf{r}$  is *stable* if

$$\lim_{p \rightarrow 0} \bar{\rho}(p) = 1, \quad (1.6)$$

<sup>1</sup>Indeed from the fact that  $\varphi_k$  are monotone and non constant for  $k \in \{1, \dots, m\}$ , it follows that  $\varphi_k(\underline{1}) = 1$ .

and *completely unstable* if  $\bar{\rho}(p) = 0$  for all  $p > 0$ . A simple coupling argument shows that  $p \mapsto \bar{\rho}(p)$  is non-increasing, so if we let

$$p_c := \sup \{p \in [0, 1] : \bar{\rho}(p) > 0\} \quad (1.7)$$

it holds that  $\bar{\rho}(p) > 0$  for all  $p < p_c$  and  $\bar{\rho}(p) = 0$  for all  $p > p_c$ . In particular, complete instability corresponds to  $p_c = 0$ .

For deterministic monotone cellular automata, Toom [Too80] has completely solved the problem of determining whether a given cellular automaton is stable or not. To state his result we first need to define *eroders*. For each local map  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$ , we let  $\Psi_\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}^{\mathbb{Z}^d}$  be defined as

$$\Psi_\varphi(x)(i) := \varphi((x(i+j))_{j \in \mathbb{Z}^d}) \quad (x \in \{0, 1\}^{\mathbb{Z}^d}), \quad (1.8)$$

i.e.,  $\Psi_\varphi$  describes one step of the time evolution of the deterministic cellular automaton defined by  $\varphi$ .

**Definition 1 (Eroders)** *We say that a local map  $\varphi$  is an eroder if for each configuration  $x \in \{0, 1\}^{\mathbb{Z}^d}$  that contains only finitely many zeros, there is a  $t \in \mathbb{N}$  such that  $\Psi_\varphi^t(x) = \underline{1}$ , where  $\Psi_\varphi^t$  denotes the  $t$ -th iterate of the map  $\Psi_\varphi$ .*

We quote the following result from [Too80, Thm 5].<sup>2</sup>

**Theorem 2 (Toom's stability theorem)** *The deterministic monotone cellular automaton  $\Phi^0$  defined by a monotone local nonconstant map  $\varphi$  is stable if  $\varphi$  is an eroder and completely unstable if  $\varphi$  is not an eroder.*

For general local maps that need not be monotone, it is known that there exists no algorithm to decide whether a given map is an eroder, even in one dimension [Pet87]. By contrast, for monotone local maps, there exists a simple criterion to check whether a given map is an eroder. To state this criterion we need the notion of minimal one-sets. A *one-set* of a monotone local map  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  is a finite set  $A \subset \mathbb{Z}^d$  such that  $\varphi(1_A) = 1$ , where  $1_A$  denotes the indicator function of  $A$ . A *minimal one-set* is a one-set that does not contain other one-sets as a proper subset. Each monotone local map  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  can be written as

$$\varphi(x) = \bigvee_{A \in \mathcal{A}(\varphi)} \bigwedge_{i \in A} x(i) \quad (x \in \{0, 1\}^{\mathbb{Z}^d}), \quad (1.9)$$

where  $\mathcal{A}(\varphi)$  is the set of minimal one-sets of  $\varphi$ . In (1.9), we use the convention that the supremum (resp. infimum) over an empty set is 0 (resp. 1). In line with this,  $\mathcal{A}(\varphi^0) = \emptyset$  and  $\mathcal{A}(\varphi^1) = \{\emptyset\}$  (note the difference!). We let  $\text{Conv}(A)$  denote the convex hull of a set  $A$ , viewed as a subset of  $\mathbb{R}^d$ . Then [Too80, Thm 6], with simplifications due to [Pon13, Thm 1], says the following.

**Proposition 3 (Erosion criterion)** *A monotone local map  $\varphi \neq \varphi^0$  is an eroder if and only if*

$$\bigcap_{A \in \mathcal{A}(\varphi)} \text{Conv}(A) = \emptyset. \quad (1.10)$$

See also Lemma 10 which gives a related alternative erosion criterion due to [Pon13, Lemma 12].

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<sup>2</sup>In the case where  $\varphi$  is not an eroder, Theorem 5 in [Too80] only states that the monotone cellular automaton is not stable, but the proof actually implies that it is completely unstable. We will give our own proof in Subsection 3.3 below.

**Remark 4** *Helly's theorem [Roc70, Corollary 21.3.2] guarantees that if (1.10) holds, then there exists a subset  $\mathcal{A}' \subset \mathcal{A}(\varphi)$  of cardinality at most  $d + 1$  such that  $\bigcap_{A \in \mathcal{A}'} \text{Conv}(A) = \emptyset$ .*

For concreteness, let us look at some examples of maps in two dimensions.

$$\begin{aligned}\varphi^{\text{NEC}}(x) &:= \text{round}([x(0,0) + x(0,1) + x(1,0)]/3), \\ \varphi^{\text{NN}}(x) &:= \text{round}([x(0,0) + x(0,1) + x(1,0) + x(0,-1) + x(-1,0)]/5), \\ \varphi^{\text{coop}}(x) &:= x(0,0) \vee (x(0,1) \wedge x(1,0)),\end{aligned}\tag{1.11}$$

where  $\text{round}$  denotes the function that rounds off a real number to the nearest integer. The function  $\varphi^{\text{NEC}}$  is known as *North-East-Center voting* or *NEC voting*, for short, and also as *Toom's rule*. In analogy with  $\varphi^{\text{NEC}}$ , we also define maps  $\varphi^{\text{NWC}}, \varphi^{\text{SWC}}, \varphi^{\text{SEC}}$  that describe North-West-Center voting, South-West-Center voting, and South-East-Center voting, respectively, defined in the obvious way. We will call the map  $\varphi^{\text{NN}}$  from (1.11) *Nearest Neighbour voting* or *NN voting*, for short. Another name found in the literature is the *symmetric majority rule*. We call  $\varphi^{\text{coop}}$  the *cooperative branching rule*. It is also known as the *sexual reproduction rule* because of the interpretation that when  $\varphi^{\text{coop}}$  is applied at a site  $(i_1, i_2)$ , two parents at  $(i_1 + 1, i_2)$  and  $(i_1, i_2 + 1)$  produce offspring at  $(i_1, i_2)$ , provided the parents' sites are both occupied and  $(i_1, i_2)$  is vacant. Using Proposition 3 one can easily check that  $\varphi^{\text{NEC}}$  and  $\varphi^{\text{coop}}$  are eroders, but  $\varphi^{\text{NN}}$  is not. Indeed, we have

$$\begin{aligned}\mathcal{A}(\varphi^{\text{NEC}}) &:= \{\{(0,0), (1,0)\}, \{(0,0), (0,1)\}, \{(0,1), (1,0)\}\}, \\ \mathcal{A}(\varphi^{\text{coop}}) &:= \{\{(0,0)\}, \{(0,1), (1,0)\}\},\end{aligned}\tag{1.12}$$

and both sets satisfy condition (1.10). On the other hand,  $\mathcal{A}(\varphi^{\text{NN}})$  is the set of all subsets of cardinality 3 of  $\{(0,0), (0,1), (1,0), (0,-1), (-1,0)\}$ . Therefore, each  $A \in \mathcal{A}(\varphi^{\text{NN}})$  contains the origin in its convex hull, and the erosion condition (1.10) is not satisfied. In fact, it is not hard to find configurations containing only finitely many zeros which cannot disappear under iterated applications of the map  $\Psi_{\varphi^{\text{NN}}}$ , for example the configuration that is zero on  $(0,0), (0,1), (1,0), (1,1)$  and one everywhere else.

**Remark 5** *Toom's stability theorem is stated in a slightly greater generality. The deterministic monotone cellular automata considered in [Too80] are defined by monotone local maps  $\varphi$  that can "look back" more than one time step, in the sense that the set  $\Delta(\varphi)$  defined above (1.1) is a finite subset of  $\mathbb{Z}^d \times \mathbb{Z}_-$ . In this case,  $\varphi$  is an eroder if and only if*

$$\bigcap_{A \in \mathcal{A}(\varphi)} \bigcup_{\alpha > 0} \{\alpha \cdot (i, t) : (i, t) \in \text{Conv}(A)\} = \emptyset,\tag{1.13}$$

that is no ray in  $\mathbb{R}^{d+1}$  that starts from the origin intersects all the convex hulls of the minimal one-sets. Note that in our setting this condition is equivalent to (1.10).

Toom's Theorem 2 settles the stability issue for deterministic monotone cellular automata. The next natural step is to study stability for monotone cellular automata with intrinsic randomness. One might think that stability should hold at least in the case when  $\varphi_1, \dots, \varphi_m$  are all eroders, but this is not true. For example, the monotone cellular automaton that applies the maps  $\varphi^{\text{NEC}}, \varphi^{\text{NWC}}, \varphi^{\text{SWC}}, \varphi^{\text{SEC}}$  each with probability 1/4 is believed to be unstable, in spite of the fact that each of these maps individually is an eroder.

To see a further example of the difficulties of cellular automata with intrinsic randomness, consider the *identity map*, defined as

$$\varphi^{\text{id}}(x) := x(0) \quad (x \in \{0,1\}^{\mathbb{Z}^d}).\tag{1.14}$$

In terms of the associated Markov chain (1.3), applying the identity map in a given space-time point has the effect that the local state at a site does not change. One might think that if  $\varphi$  is an eroder, then a cellular automaton that applies the maps  $\varphi$  and  $\varphi^{\text{id}}$  each with positive probability must be stable, but again this turns out to be wrong. Gray [Gra99, Examples 18.3.5 and 18.3.6] has given convincing arguments that show that the addition of the identity map can make eroders unstable and conversely, make non-eroders stable. Being able to include the identity map is important for understanding continuous-time interacting particle systems. We can think of such systems as limits of discrete-time cellular automata where time is measured in steps of some small size  $\delta$  and all maps except  $\varphi^{\text{id}}$  are applied with a probability of order  $\delta$ .

The most difficult part of Theorem 2 is the statement that  $\Phi^0$  is stable if  $\varphi$  is an eroder. To prove this, Toom used an intricate Peierls argument. It is fair to say that Toom’s original paper [Too80] is quite hard to read. Indeed, several subsequent papers have been devoted to simplifying his arguments and others have re-proved his result from scratch for some specific model at hand to avoid relying on this complex proof [LMS90, Gac95, Pre07, Pon13, Gac21] (see Subsection 1.5).

In this paper, we reformulate and simplify Toom’s Peierls argument. Our main motivation is the problem of extending Toom’s stability theorem to monotone cellular automata with intrinsic randomness. As a first step in this direction, we will prove a stability result in Theorem 9 below, which however excludes many interesting cases such as cellular automata that apply the identity map with a positive probability. This is not due to a fundamental limitation of Toom’s Peierls argument, but to go beyond Theorem 9 one needs more advanced methods to estimate the Peierls sum. In order not to overload the present paper, we have delegated these methods to a companion paper [SST24] where further stability results for cellular automata with intrinsic randomness will be proved.

As a further result of our reformulation of Toom’s Peierls argument, we will derive explicit lower bounds on the critical noise parameter  $p_c$  from (1.7) for some deterministic cellular automata. Although these bounds are often several orders of magnitude from the conjectured true values, they are nevertheless the sharpest rigorous bounds available. For a subclass of cellular automata, we show that it is possible to derive significantly better bounds by using Toom cycles, which are Toom contours with additional pleasant properties.

Toom’s Peierls argument was invented to study stability of monotone cellular automata with respect to noise that is i.i.d. in space and time. It has recently been discovered that it can also be used to prove stability with respect to noise that is applied only to the initial state [HS22, CSS24]. This has applications in bootstrap percolation, which we will briefly discuss in Subsection 1.4 below.

## Outline

In the remainder of Section 1 we discuss applications of Toom’s Peierls argument. Stability of monotone cellular automata with intrinsic randomness is discussed in Subsection 1.2, explicit lower bounds on the critical noise parameter are presented in Subsection 1.3, and bootstrap percolation is discussed in Subsection 1.4. In Subsection 1.5 we discuss earlier work on the topic and state some open problems.

In Section 2 we present our reformulation of Toom’s Peierls argument. We work in a more general setting than in Section 1, which also allows for cellular automata that can look back more than one time step as in Remark 5 and cellular automata on other lattices than  $\mathbb{Z}^d$ , such as trees. We show that each monotone cellular automaton has a maximal trajectory and that the density  $\bar{p}(p)$  of the upper invariant law is equal to the probability that this maximal trajectory has a one at the origin. We moreover introduce objects we call *Toom contours* that are directed graphs with different types of edges that are designed to make use of the

characterisation of eroders in terms of edge speeds and polar functions that will be discussed in Subsection 1.2 below.

The main results of Section 2 and indeed of the whole paper are Theorems 23, 26, and 27. Theorem 23 says that on the event that the maximal trajectory has a zero at the origin, a Toom contour must be present. As stated precisely in Theorem 27, this allows one to estimate  $1 - \bar{\rho}(p)$  from above by the expected number of Toom contours that are present in a cellular automaton. This is the core of Toom's Peierls argument. Theorem 26 shows that for a subclass of monotone cellular automata, it is possible to work with *Toom cycles* which are Toom contours with additional pleasant properties that often lead to sharper bounds.

The remainder of the paper is devoted to proofs. In Section 3 we prove some preparatory results, in Section 4 we prove the results from Section 2 about Toom contours, and in Section 5 we apply these results to prove stability of a class of monotone cellular automata with intrinsic randomness and derive some explicit bounds on the critical noise parameter.

## 1.2 Stability of monotone cellular automata

In this subsection we state a theorem giving sufficient conditions for the stability of monotone cellular automata with intrinsic randomness. The statement of the theorem involves edge speeds and gives additional insight into Toom's stability theorem, which it generalises.

**Definition 6 (Edge speed)** *Let  $\ell : \mathbb{R}^d \rightarrow \mathbb{R}$  be a linear form that is not identically zero and let  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  be a monotone local map. We call the quantity*

$$\varepsilon_\varphi(\ell) := \sup_{A \in \mathcal{A}(\varphi)} \inf_{i \in A} \ell(i). \quad (1.15)$$

*the edge speed of  $\varphi$  in the direction  $\ell$ .*

The name ‘‘edge speed’’ already suggests its interpretation. For any linear form  $\ell : \mathbb{R}^d \rightarrow \mathbb{R}$ , let  $H_r^\ell \in \{0, 1\}^{\mathbb{Z}^d}$  denote the half-space configuration defined by

$$H_r^\ell(i) := \begin{cases} 1 & \text{if } \ell(i) \geq r, \\ 0 & \text{if } \ell(i) < r \end{cases} \quad (r \in \mathbb{R}). \quad (1.16)$$

The following lemma explains the name ‘‘edge speed’’.

**Lemma 7 (Edge speeds)** *Let  $\ell : \mathbb{R}^d \rightarrow \mathbb{R}$  be a linear form that is not identically equal to zero and  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  be a monotone local map. Then for each  $r \in \mathbb{R}$  and  $t \geq 0$  the map from (1.8) satisfies*

$$\Psi_\varphi^t(H_r^\ell) = H_{r-t\varepsilon_\varphi(\ell)}^\ell. \quad (1.17)$$

The result above, which follows easily from the definitions, is for completeness proved in Subsection 3.1. To state our stability result, we need one more definition.

**Definition 8 (Polar functions)** *Given an integer  $\sigma \geq 2$ , a polar function of dimension  $\sigma$  is a linear function*

$$\mathbb{R}^d \ni z \mapsto L(z) = (L_1(z), \dots, L_\sigma(z)) \in \mathbb{R}^\sigma \quad (1.18)$$

*such that*

$$\sum_{s=1}^{\sigma} L_s(z) = 0 \quad (z \in \mathbb{R}^d). \quad (1.19)$$

In Subsection 5.2 we will use Toom contours to prove the following stability result.

**Theorem 9 (Stability of monotone cellular automata with intrinsic randomness)**  
Fix  $m \geq 1$  and let  $\Phi^0$  be a monotone cellular automaton defined by maps  $\varphi_1, \dots, \varphi_m$  and a probability distribution  $\mathbf{r}(1), \dots, \mathbf{r}(m)$ . Assume that there exists a linear polar function  $L$  of dimension  $\sigma \geq 2$  such that the worst-case edge speeds

$$\varepsilon_s := \inf_{1 \leq k \leq m} \varepsilon_{\varphi_k}(L_s) \quad (1 \leq s \leq \sigma) \quad (1.20)$$

satisfy

$$\varepsilon := \sum_{s=1}^{\sigma} \varepsilon_s > 0. \quad (1.21)$$

Then  $\Phi^0$  is stable.

Theorem 9 is far from optimal in terms of what can be achieved by Toom's Peierls argument, but to improve on it one needs more advanced methods to estimate the Peierls sum which will be presented in our companion paper [SST24]. Although Theorem 9 is suboptimal in the presence of intrinsic randomness, it is optimal in the deterministic case. To see this, we need the following alternative erosion criterion originally due to [Pon13, Lemma 12], the proof of which will be given in Subsection 3.1.

**Lemma 10 (Alternative erosion criterion)** Let  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  be a non-constant monotone function. Then  $\varphi$  is an eroder if and only if there exists a polar function  $L$  of dimension  $\sigma \geq 2$  such that the edge speeds defined in (1.15) satisfy

$$\sum_{s=1}^{\sigma} \varepsilon_{\varphi}(L_s) > 0. \quad (1.22)$$

It is instructive to see why (1.22) implies that  $\varphi$  is an eroder. Given a configuration  $x \in \{0, 1\}^{\mathbb{Z}^d}$  containing finitely many zeros, let the *extent* of  $x$  be defined as

$$\text{ext}(x) := \begin{cases} \sum_{s=1}^{\sigma} r_s(x) & \text{with } r_s(x) := \sup \{L_s(i) : i \in \mathbb{Z}^d, x(i) = 0\} & \text{if } x \neq \underline{1}, \\ -\infty & & \text{if } x = \underline{1} \end{cases} \quad (1.23)$$

By the defining property (1.19) of a linear polar function,  $\text{ext}(x) \geq 0$  for each  $x \neq \underline{1}$ . Lemma 7 and the monotonicity of  $\varphi$  imply that for each configuration  $x$  with finitely many zeros,

$$\text{ext}(\Psi_{\varphi}^t(x)) \leq \text{ext}(x) - \varepsilon_{\varphi}(L)t \quad (t \geq 0), \quad (1.24)$$

and hence  $\Psi_{\varphi}^t(x) = \underline{1}$  for all  $t \geq \text{ext}(x)/\varepsilon$ . In the case with intrinsic randomness, condition (1.21) similarly implies that if  $(X_t)_{t \geq 0}$  is the Markov chain defined as in (1.3) in terms of the unperturbed automaton  $\Phi^0$ , started in an initial state  $X_0 = x$  with finitely many zeros, then almost surely

$$\text{ext}(X_t) \leq \text{ext}(x) - \varepsilon t \quad (t \geq 0), \quad (1.25)$$

and  $X_t = \underline{1}$  for all  $t \geq \text{ext}(x)/\varepsilon$ . Thus Theorem 9 proves stability under the assumption that under the unperturbed evolution, finite collections of zeros disappear after a finite *deterministic* time. There are many examples of monotone cellular automata with intrinsic randomness that do not satisfy (1.21) but for which under the unperturbed evolution, finite collections of zeros disappear after a finite *random* time. For some of these, we will prove stability in our companion paper [SST24].



### 1.3 Bounds on the critical noise parameter

In this subsection we apply Theorem 9 to some concrete examples and derive explicit bounds on the critical noise parameter  $p_c$  from (1.7).

We first set  $m = 1$  and consider the deterministic cellular automaton on  $\mathbb{Z}^2$  defined by the single map  $\varphi_1 = \varphi^{\text{coop}}$ . The function  $L : \mathbb{R}^2 \rightarrow \mathbb{R}^2$  defined as

$$L_1(z) := -z_1 - z_2, \quad L_2(z) := z_1 + z_2 \quad (z = (z_1, z_2) \in \mathbb{R}^2) \quad (1.26)$$

is a linear polar function of dimension  $\sigma = 2$  in the sense of Definition 8. Using (1.12) and (1.15), we see that the corresponding edge speeds from (1.20) are given by

$$\varepsilon_1 = \varepsilon_{\varphi^{\text{coop}}}(L_1) = 0, \quad \varepsilon_2 = \varepsilon_{\varphi^{\text{coop}}}(L_2) = 1, \quad (1.27)$$

so  $\varepsilon = \varepsilon_1 + \varepsilon_2 > 0$  and hence Theorem 9 implies that this cellular automaton is stable.

We next set  $m = 1$  and  $\varphi_1 = \varphi^{\text{NEC}}$ . We define a linear polar function  $L$  of dimension  $\sigma = 3$  by

$$L_1(z_1, z_2) := -z_1, \quad L_2(z_1, z_2) := -z_2, \quad L_3(z_1, z_2) := z_1 + z_2 \quad (z \in \mathbb{R}^2). \quad (1.28)$$

One can check that for this choice of  $L$  (recall (1.12))

$$\varepsilon_1 = \varepsilon_{\varphi^{\text{NEC}}}(L_1) = 0, \quad \varepsilon_2 = \varepsilon_{\varphi^{\text{NEC}}}(L_2) = 0, \quad \varepsilon_3 = \varepsilon_{\varphi^{\text{NEC}}}(L_3) = 1, \quad (1.29)$$

which implies  $\varepsilon = \varepsilon_1 + \varepsilon_2 + \varepsilon_3 > 0$ , hence stability.

To also see an example with intrinsic randomness, consider the case  $m = 2$  with  $\varphi_1 = \varphi^{\text{NEC}}$  and  $\varphi_2 = \varphi^{\text{coop}}$ . Using the polar function (1.28) one can check that

$$\varepsilon_1 = \varepsilon_{\varphi^{\text{NEC}}}(L_1) \wedge \varepsilon_{\varphi^{\text{coop}}}(L_1) = 0 \wedge 0 = 0, \quad (1.30)$$

$$\varepsilon_2 = \varepsilon_{\varphi^{\text{NEC}}}(L_2) \wedge \varepsilon_{\varphi^{\text{coop}}}(L_2) = 0 \wedge 0 = 0, \quad (1.31)$$

$$\varepsilon_3 = \varepsilon_{\varphi^{\text{NEC}}}(L_3) \wedge \varepsilon_{\varphi^{\text{coop}}}(L_3) = 1 \wedge 1 = 1, \quad (1.32)$$

which implies  $\varepsilon > 0$ . Therefore, Theorem 9 implies stability for this cellular automaton regardless of the choice of the probability distribution  $\mathbf{r} = (\mathbf{r}(1), \mathbf{r}(2))$  on  $\{1, 2\}$ .

To see an example where Theorem 9 is not applicable, consider the case  $m = 4$  with  $\varphi_1 = \varphi^{\text{NEC}}$ ,  $\varphi_2 = \varphi^{\text{NWC}}$ ,  $\varphi_3 = \varphi^{\text{SWC}}$ , and  $\varphi_4 = \varphi^{\text{SEC}}$ . In this case, there exists no polar function that satisfies the hypothesis of Theorem 9. We conjecture that this model is unstable if  $\mathbf{r}$  is the uniform distribution on  $\{1, 2, 3, 4\}$ , but stable in all other cases.

The proof of Theorem 9 allows us to derive explicit lower bounds on the critical noise parameter  $p_c$  from (1.7). In particular, in Subsection 5.7 we will prove the following bounds.

**Proposition 11 (Explicit bounds)** *For the deterministic cellular automaton on  $\mathbb{Z}^2$  that applies  $\varphi^{\text{coop}}$  in each space-time point  $p_c \geq 1/64$ . For the deterministic cellular automaton on  $\mathbb{Z}^2$  that applies  $\varphi^{\text{NEC}}$  in each space-time point  $p_c \geq 3^{-21}$ .*

In our companion paper [SST24], using a more advanced method to bound the Peierls sum, we will improve the lower bound for the cellular automaton defined by  $\varphi^{\text{NEC}}$  to  $p_c \geq 1/12000$ . Numerical simulations suggest that the true value of  $p_c$  is  $\approx 0.105$  for  $\varphi^{\text{coop}}$  and  $\approx 0.053$  for  $\varphi^{\text{NEC}}$ . There is a good reason why the rigorous bounds for  $\varphi^{\text{NEC}}$  are worse than for  $\varphi^{\text{coop}}$ . If we want to apply Lemma 10 to prove that  $\varphi^{\text{NEC}}$  is an eroder, then we need a linear polar function of dimension at least three, while for  $\varphi^{\text{coop}}$  a linear polar function of dimension two suffices. In general, the higher the dimension of the linear polar function, the worse the bounds. For linear polar functions of dimension two, we can moreover use Toom cycles instead of Toom contours, which also leads to sharper bounds.

## 1.4 Bootstrap percolation

Our simplification of Toom's argument has successfully been used in the study of bootstrap percolation in [HS22]. In this subsection, we briefly elaborate on this connection.

Instead of perturbing a deterministic monotone cellular automaton with noise that is i.i.d. in space and time, one can also look at perturbations that are i.i.d. in space but constant in time. More precisely, if  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  is a monotone local map that is not constant, then it is interesting to look at i.i.d. collections of random variables  $\Phi^p = (\Phi_i^p)_{i \in \mathbb{Z}^d}$  with values in  $\{\varphi^0, \varphi\}$  such that

$$\mathbb{P}[\Phi_i^p = \varphi^0] = p \quad \text{and} \quad \mathbb{P}[\Phi_i^p = \varphi] = 1 - p, \quad (1.33)$$

and modify the evolution in (1.3) in the sense that the same map  $\Phi_i^p$  is applied at each time  $t > 0$ . We let  $\tilde{\rho}_\varphi(p)$  denote the long-time limit of the density started from  $\underline{1}$  for this type of evolution. Recall from (1.5) that  $\bar{\rho}_\varphi(p)$  denotes the limiting density when the noise is i.i.d. in space and time.

**Definition 12 (Forms of stability)** *We say that  $\varphi$  is stable in the bootstrap sense if  $\lim_{p \rightarrow 0} \tilde{\rho}_\varphi(p) = 1$  and stable in Toom's sense if  $\lim_{p \rightarrow 0} \bar{\rho}_\varphi(p) = 1$ .*

Toom's stability theorem (Theorem 2) completely answers the question which monotone local maps  $\varphi$  are stable in Toom's sense. The analogue question for stability in the bootstrap sense has been answered more recently. In order to formulate a precise result, we first need to translate the problem into the language of bootstrap percolation. In analogy with notation introduced in (1.8), we let  $\Psi_\varphi^p : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}^{\mathbb{Z}^d}$  denote the random map defined as

$$\Psi_\varphi^p(x)(i) := \Phi_i((x(i+j))_{j \in \mathbb{Z}^d}) \quad (x \in \{0, 1\}^{\mathbb{Z}^d}). \quad (1.34)$$

Since the noise is the same in each time step,

$$\tilde{\rho}_\varphi(p) := \lim_{t \rightarrow \infty} \mathbb{P}[(\Psi_\varphi^p)^t(\underline{1})(i) = 1] \quad (i \in \mathbb{Z}^d), \quad (1.35)$$

where  $(\Psi_\varphi^p)^t$  denotes the  $t$ -th iterate of the map  $\Psi_\varphi^p$ . Note that if  $(\Psi_\varphi^p)^t(\underline{1})(i) = 0$  for some  $i \in \mathbb{Z}^d$  and  $t \geq 1$ , then  $(\Psi_\varphi^p)^{t+s}(\underline{1})(i) = 0$  for all  $s \geq 0$ .

*Minimal zero-sets* of a monotone local map  $\varphi$  are defined analogously to the minimal one-sets of (1.9), i.e., these are minimal elements of the set of all finite  $Z \subset \mathbb{Z}^d$  with the property that  $\varphi(1 - 1_Z) = 0$ . We let  $\mathcal{Z}(\varphi)$  denote the set of all minimal zero-sets of  $\varphi$ . For any monotone local map  $\varphi$  we define a map  $\bar{\varphi} : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  in terms of its set of minimal zero-sets as

$$\mathcal{Z}(\bar{\varphi}) := \{Z : Z \in \mathcal{Z}(\varphi), 0 \notin Z\}. \quad (1.36)$$

We say that two monotone local maps  $\varphi_1$  and  $\varphi_2$  from  $\{0, 1\}^{\mathbb{Z}^d}$  to  $\{0, 1\}$  are *equivalent in the bootstrap sense* if  $\bar{\varphi}_1 = \bar{\varphi}_2$ . Observe that in this case  $\varphi_1(x) = \varphi_2(x)$  for all  $x \in \{0, 1\}^{\mathbb{Z}^d}$  such that  $x(0) = 1$ . It is easy to see that if  $\varphi_1$  and  $\varphi_2$  are equivalent in the bootstrap sense, then  $(\Psi_{\varphi_1}^p)^t(\underline{1})$  and  $(\Psi_{\varphi_2}^p)^t(\underline{1})$  are equal in law, and moreover almost surely equal if we couple both processes in the obvious way, by applying the zero map in the same space-time points. Therefore, in view of (1.35), for equivalent maps  $\tilde{\rho}_{\varphi_1}(p) = \tilde{\rho}_{\varphi_2}(p)$  ( $0 \leq p \leq 1$ ). In particular, for each map  $\varphi$

$$\tilde{\rho}_\varphi(p) = \tilde{\rho}_{\bar{\varphi}}(p). \quad (1.37)$$

The set  $\mathcal{Z}(\bar{\varphi})$  is called the *update family* in the bootstrap percolation literature, and it is traditionally in terms of this set that the dynamics are described.

Bootstrap percolation was first introduced in [CLR79] to model magnetic materials at low temperature, and has since been extensively studied (see [Mor17] for a review). The first stability result on  $\mathbb{Z}^d$  was established in [Sch92] for the update family consisting of all sets

containing exactly  $r$  neighbours of the origin. Presently, there is a complete characterisation of bootstrap percolation maps on  $\mathbb{Z}^d$ . Based on the geometry of the sets in the update family, bootstrap maps  $\varphi$  fall into three universality classes: maps  $\varphi$  in the supercritical and critical classes are completely unstable in the bootstrap sense meaning that  $\tilde{\rho}_\varphi(p) = 0$  for all  $p > 0$ , while those in the subcritical class are stable in the bootstrap sense. These results were first established in [BSU15, BB+16] in two dimensions and recently extended to higher dimensions in a series of papers [BB+22a, BB+22b, BB+24]. It was shown in [HS22] that the stability result in the subcritical class can also (and more simply) be obtained using Toom contours, marking the first application of the results presented in this paper. Additionally, [HT24] applies Toom contours to establish exponential decay results for the largest cluster of zeros in the final configuration of subcritical bootstrap percolation, as well as for the largest cluster of zeros in space-time for monotone cellular automata defined by an eroder.

The following theorem completely answers the question which monotone local maps  $\varphi$  are stable in the bootstrap sense.

**Theorem 13 (Stability in the bootstrap sense)** *Let  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  be a monotone local map. Then  $\varphi$  is stable in the bootstrap percolation sense if and only if there exists a linear polar function  $L$  of dimension  $\sigma \geq 2$  such that the edge speeds defined in (1.15) of the map  $\bar{\varphi}$  from (1.36) satisfy*

$$\varepsilon_{\bar{\varphi}}(L_s) > 0 \quad (1 \leq s \leq \sigma). \quad (1.38)$$

This theorem is proved in [HS22, Lemma 2.1]. (The lemma only states one direction of the theorem, but it is easy to check that every implication used in the proof is in fact an equivalence.) We observe that (1.38) is stronger than the condition that  $\sum_{s=1}^{\sigma} \varepsilon_{\bar{\varphi}}(L_s) > 0$ , so comparing with Theorem 2 and Lemma 10 we see that stability in the bootstrap sense implies that  $\bar{\varphi}$  is stable in Toom's sense but not vice versa. To see a concrete example, consider the Duarte model on  $\mathbb{Z}^2$  defined by the monotone local map

$$\varphi^{\text{Duarte}}(x) := \text{round}([x(0, 1) + x(0, -1) + x(-1, 0)]/3). \quad (1.39)$$

This rule is unstable in the bootstrap sense as it is known to belong to the critical universality class. However, as the intersection of the convex hulls of its minimal one-sets is empty,  $\varphi^{\text{Duarte}}$  is an eroder and hence stable in Toom's sense. For any monotone local map  $\varphi$  we can apart from the map  $\bar{\varphi}$  from (1.36) also define

$$\underline{\mathcal{Z}}(\varphi) := \underline{\mathcal{Z}}(\bar{\varphi}) \cup \{0\}.$$

Then  $\underline{\varphi}$  is equivalent to  $\varphi$  in the bootstrap sense<sup>3</sup> while it is easy to see that  $\underline{\varphi}$  is never stable in Toom's sense, as once a site flips to 0 it forever remains in state 0. Thus, for general monotone local maps  $\varphi$  all four combinations (stable/unstable in Toom's/the bootstrap sense) are possible.

## 1.5 Discussion

The cellular automaton defined by the NEC voting map  $\varphi^{\text{NEC}}$  is nowadays known as *Toom's model*. In line with Stigler's law of eponymy, Toom's model was not invented by Toom, but by Vasilyev, Petrovskaya, and Pyatetski-Shapiro, who simulated random perturbations of this and other models on a computer [VPP69]. Toom, having heard of [VPP69] during a seminar, proved in [Too74] that there exist random cellular automata on  $\mathbb{Z}^d$  with at least  $d$  different invariant laws. Although Toom's model is not explicitly mentioned in the paper, his proof method can be applied to prove that  $p_c > 0$  for his model. In [Too80], Toom improved his methods and proved his celebrated stability theorem. His paper is quite hard to read. A more

<sup>3</sup>In fact, it is easy to see that  $\underline{\varphi}$  and  $\bar{\varphi}$  are the smallest and largest maps that are equivalent to  $\varphi$ .

accessible account of Toom’s original argument (with pictures!) in the special case of Toom’s model can be found in the appendix of [LMS90].<sup>4</sup>

Bramson and Gray [BG91] have given another alternative proof of Toom’s stability theorem that relies on comparison with continuum models (which describe unions of convex sets in  $\mathbb{R}^d$  evolving in continuous time) and renormalisation-style block arguments. A disadvantage of this approach is that it is restricted to lattices that can be rescaled to  $\mathbb{R}^d$  while Toom’s method can also work on lattices such as trees, as demonstrated in [CSS24]. Gray [Gra99] proved a stability theorem for monotone interacting particle systems (i.e., in continuous time). The proofs use ideas from [Too80] and [BG91]. Gray also derived necessary and sufficient conditions for a monotonic map to be an eroder [Gra99, Thm 18.2.1], apparently overlooking the fact that Toom had already proved the much simpler condition (1.10).

The cellular automaton that applies the monotone map  $\varphi$  with probability  $p$  and the identity map  $\varphi^{\text{id}}$  with probability  $1 - p$  is also referred to in the literature as *p-asynchronous* cellular automaton. In asynchronous cellular automata, cells do not update their states simultaneously. There are various ways to define this asynchrony; for a comprehensive survey, see [Fat13]. In [Gha92], a generalization of Toom’s theorem was presented for a particular class of asynchronous cellular automata.

Motivated by abstract problems in computer science, a number of authors have given alternative proofs of Toom’s stability theorem in a more restrictive setting [GR88, BS88, Gac95, Gac21]. Their main interest is in a three-dimensional system which evolves in two steps: letting  $e_1, e_2, e_3$  denote the basis vectors in  $\mathbb{Z}^3$ , they first replace  $X_n(i)$  by

$$X'_n(i) := \text{round}((X_n(i) + X_n(i + e_1) + X_n(i + e_2))/3),$$

and then set

$$X_{n+1}(i) := \text{round}((X'_n(i) + X'_n(i + e_3) + X'_n(i - e_3))/3).$$

They prove explicit bounds for finite systems, although for values of  $p$  that are extremely close to zero.<sup>5</sup> The proofs of [GR88] do not use Toom’s Peierls argument but rely on different methods. Their bounds were improved in [BS88]. Still better bounds can be found in the unpublished note [Gac95]. The proofs in the latter manuscript are very similar to Toom’s argument, with some crucial improvements at the end that are hard to follow due to missing definitions (which might explain why this manuscript remained unpublished). This version of the argument seems to have inspired the incomplete note by John Preskill [Pre07] who links it to the interesting idea of counting “minimal explanations”. We will use this general idea in Subsection 4.1 below, but our precise definition of a “minimal explanation” differs a bit from his. As explained at Figure 3 and at the end of Subsection 2.3, the relation between Toom contours and minimal explanations is not so straightforward as suggested in [Gac95, Pre07].

Around 1985, Durrett and Gray submitted a very interesting paper about an interacting particle system based on the map  $\varphi^{\text{coop}}$  from (1.11). The major revision requested by the referee never materialised, however. For many years, a short note announcing the results without proofs [Dur86] was the only accessible source to this material but recently Rick Durrett has made the original preprint available on his homepage [DG85]. Hwa-Nien Chen [Che92, Che94], who was a PhD student of Lawrence Gray, studied the stability of various variations of Toom’s model under perturbations of the initial state and the birth rate. The proofs of two of his four theorems depend on results that he cites from the preprint [DG85]. Ponselet [Pon13] gave an excellent account of the existing literature and together with her supervisor proved exponential decay of correlations for the upper invariant law of a large class of randomly perturbed monotone cellular automata [MP11].

<sup>4</sup>Unfortunately, their Figure 6 contains a small mistake, in the form of an arrow that should not be there.

<sup>5</sup>In particular, [Gac95] needs  $p < 2^{-21}3^{-8}$ .

There exists duality theory for general monotone interacting particle systems [Gra86, SS18, LS23]. The basic idea is that the state in the origin at time zero is a monotone function of the state at time  $-t$ , and this monotone function evolves in a Markovian way as a function of  $t$ . As noted in [Dur86] this dual process plays an important ingredient of the proofs of [DG85]. It is also closely related to the minimal explanations of Preskill [Pre07]. A good understanding of this dual process could potentially help solve many open problems in the area, but its behaviour is already quite complicated in the mean-field case [MSS20].

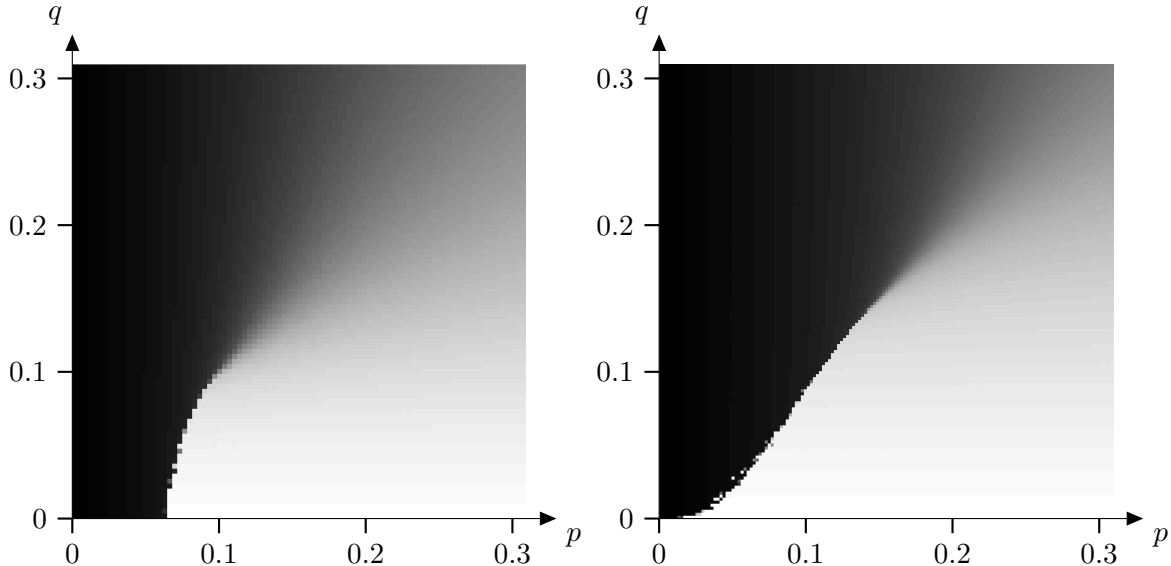


Figure 1: Density  $\bar{\rho}$  of the upper invariant law of two monotone random cellular automata as a function of the parameters, shown on a scale from 0 (white) to 1 (black). On the left: a version of Toom’s model that applies the maps  $\varphi^0$ ,  $\varphi^1$ , and  $\varphi^{\text{NEC}}$  with probabilities  $p$ ,  $q$ , and  $1 - p - q$ , respectively. On the right: the monotone random cellular automaton that applies the maps  $\varphi^0$ ,  $\varphi^1$ , and  $\varphi^{\text{NN}}$  with probabilities  $p$ ,  $q$ , and  $1 - p - q$ , respectively. The map  $\varphi^{\text{NEC}}$  is an eroder but  $\varphi^{\text{NN}}$  is not. By the symmetry between the 0’s and the 1’s, in both models, the density  $\rho(p, q)$  of the lower invariant law equals  $1 - \bar{\rho}(q, p)$ . Due to metastability effects, the area where the upper invariant law differs from the lower invariant law is shown too large in these numerical data. For Toom’s model with  $q = 0$ , the data shown above suggest a first order phase transition at  $p_c \approx 0.057$  but based on numerical data for edge speeds we believe the true value is  $p_c \approx 0.053$ . We conjecture that the model on the right has a unique invariant law everywhere except on the diagonal  $p = q$  for  $p$  sufficiently small.

In numerical investigations of monotone cellular automata, it is often useful to take a wider view and perturb the system not only with i.i.d. zeros but also with i.i.d. ones. Recall the constant maps  $\varphi^0$  and  $\varphi^1$  defined in (1.1). If  $\varphi : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  is a monotone local map that is not constant, then it is interesting to look at i.i.d. collections of random variables  $\Phi^{p,q} = (\Phi_{i,t}^{p,q})_{(i,t) \in \mathbb{Z}^{d+1}}$  with values in  $\{\varphi^0, \varphi^1, \varphi\}$  such that

$$\mathbb{P}[\Phi_{i,t}^{p,q} = \varphi^0] = p, \quad \mathbb{P}[\Phi_{i,t}^{p,q} = \varphi^1] = q, \quad \mathbb{P}[\Phi_{i,t}^{p,q} = \varphi] = 1 - p - q, \quad (1.40)$$

where  $p, q \geq 0$  with  $p + q \leq 1$  are parameters of the model. Figure 1 shows numerical data for the density  $\bar{\rho}(p, q)$  of the upper invariant law for such a cellular automaton in the case when  $\varphi = \varphi^{\text{NEC}}$  and  $\varphi = \varphi^{\text{NN}}$ , respectively. We see that in line with Toom’s theorem,  $\lim_{p \rightarrow 0} \bar{\rho}(p, 0) = 1$  for the NEC voting rule but  $\bar{\rho}(p, 0) = 0$  for all  $p > 0$  in the case of NN voting. Nevertheless, the simulations suggest that the NN voting rule is stable under

symmetric noise, in the sense that  $\lim_{p \rightarrow 0} \bar{\rho}(p, p) = 1$ . Proving this is a long-standing open problem; a continuous-time version of this model is mentioned in [Lig85, Example I.4.3(e)]. A closely related problem, that is also open, is to show that for the NEC voting rule the function  $p \mapsto \rho(p)$  makes a jump at  $p_c := \sup\{p : \bar{\rho}(p) > 0\}$ .

## 2 Toom contours

### 2.1 Monotone cellular automata

In this section, we introduce Toom contours, which are the central object in Toom's Peierls argument. Toom contours can be defined for monotone cellular automata in which space-time has a more general structure than  $\mathbb{Z}^{d+1}$ . In the present subsection, we extend the definitions of Section 1 to this more general set-up.

Let  $\Lambda$  be a countable set. As in Subsection 1.1, we say that a map  $\phi : \{0, 1\}^\Lambda \rightarrow \{0, 1\}$  is *local* if there exists a finite  $\Delta \subset \Lambda$  such that  $\phi(x)$  depends only on  $(x(i))_{i \in \Delta}$  and we let  $\Delta(\phi)$  denote the smallest such set. In analogy with (1.1), in the present setting, we denote the constant functions by

$$\phi^0(x) := 0 \quad \text{and} \quad \phi^1(x) := 1 \quad (x \in \{0, 1\}^\Lambda). \quad (2.1)$$

We let  $\mathcal{A}(\phi)$  denote the set of minimal one-sets of  $\phi$ , defined as in Subsection 1.1.

Recall that a *directed graph* is a pair  $(V, \vec{E})$  where  $V$  is a set whose elements are called *vertices* and  $\vec{E}$  is a subset of  $V \times V$  whose elements are called *directed edges*. For each directed edge  $(v, w) \in \vec{E}$ , we call  $v$  the starting vertex and  $w$  the endvertex. We say that  $(V, \vec{E})$  is *acyclic* if there do not exist  $n \geq 1$  and  $v_0, \dots, v_n \in V$  with  $v_n = v_0$  such that  $(v_{k-1}, v_k) \in \vec{E}$  for all  $0 < k \leq n$ .

Let  $\phi = (\phi_i)_{i \in \Lambda}$  be a collection of local maps  $\phi_i : \{0, 1\}^\Lambda \rightarrow \{0, 1\}$ , and let

$$\vec{H}(\phi) := \{(i, j) \in \Lambda^2 : j \in \Delta(\phi_i)\}. \quad (2.2)$$

Then  $(\Lambda, \vec{H}(\phi))$  is a directed graph. Generalising our earlier definition, we define a *cellular automaton* to be a collection of local maps  $\phi = (\phi_i)_{i \in \Lambda}$  for which the directed graph  $(\Lambda, \vec{H}(\phi))$  is acyclic. We call  $(\Lambda, \vec{H}(\phi))$  the *dependence graph* associated with  $\phi$ . A *trajectory* of a cellular automaton is a function  $x : \Lambda \rightarrow \{0, 1\}$  such that

$$x(i) = \phi_i(x) \quad (i \in \Lambda). \quad (2.3)$$

A cellular automaton is *monotone* if  $\phi_i$  is a monotone map for each  $i \in \Lambda$ , i.e.,  $x \leq y$  (coordinatewise) implies  $\phi_i(x) \leq \phi_i(y)$ .

To make the link with our earlier definitions from Subsection 1.1, let  $\Phi^p = (\Phi_{i,t}^p)_{(i,t) \in \mathbb{Z}^{d+1}}$  be a monotone cellular automaton of the type considered in (1.2), and for each  $(i, t) \in \mathbb{Z}^{d+1}$ , define  $\Phi_{(i,t)}^p : \{0, 1\}^{\mathbb{Z}^{d+1}} \rightarrow \{0, 1\}$  by

$$\Phi_{(i,t)}^p((x(i', t'))_{(i', t') \in \mathbb{Z}^{d+1}}) := \Phi_{i,t}^p((x(i + i', t - 1))_{i' \in \mathbb{Z}^d}) \quad (x \in \{0, 1\}^{\mathbb{Z}^{d+1}}). \quad (2.4)$$

Then  $(\Phi_{(i,t)}^p)_{(i,t) \in \mathbb{Z}^{d+1}}$  is a random monotone cellular automaton according to the definitions of the present section. Note the subtle difference in notation between  $\Phi_{i,t}^p$  and  $\Phi_{(i,t)}^p$ . By a slight abuse of notation, we use the symbol  $\Phi^p$  for both the collections  $(\Phi_{i,t}^p)_{(i,t) \in \mathbb{Z}^{d+1}}$  and  $(\Phi_{(i,t)}^p)_{(i,t) \in \mathbb{Z}^{d+1}}$ .

We next turn our attention to the lower and upper invariant laws from formula (1.4). The following two lemmas introduce two closely related objects, the minimal and maximal trajectories, and show how they are related to the lower and upper invariant laws. We prove these lemmas in Subsection 3.2.

**Lemma 14 (Minimal and maximal trajectories)** *Let  $\phi$  be a monotone cellular automaton. Then there exist trajectories  $\underline{x}$  and  $\bar{x}$  that are uniquely characterised by the property that each trajectory  $x$  of  $\phi$  satisfies  $\underline{x} \leq x \leq \bar{x}$  (pointwise).*

**Lemma 15 (Lower and upper invariant laws)** *Let  $\Phi^p$  be the random monotone cellular automaton defined in (2.4) and let  $\underline{X}^p$  and  $\bar{X}^p$  be its minimal and maximal trajectories. Then*

$$\mathbb{P}[(\underline{X}^p(i, t))_{i \in \mathbb{Z}^d} \in \cdot] = \underline{\nu}_p \quad \text{and} \quad \mathbb{P}[(\bar{X}^p(i, t))_{i \in \mathbb{Z}^d} \in \cdot] = \bar{\nu}_p \quad (t \in \mathbb{Z}), \quad (2.5)$$

where  $\underline{\nu}_p$  and  $\bar{\nu}_p$  are the lower and upper invariant laws of the Markov chain in (1.3).

Let  $\Phi$  be a random monotone cellular automaton, i.e., a random variable taking values in the space of all monotone cellular automata on a given space-time set  $\Lambda$ , and let  $\bar{X}$  denote its maximal trajectory, which is now also random. In Theorem 27 below, we give a lower bound on the probability  $\mathbb{P}[\bar{X}(i) = 1]$ . We will show that on the event that  $\bar{X}(i) = 0$ , the random monotone cellular automaton  $\Phi$  must contain a certain structure that we will call a *Toom contour rooted at  $i$* . The probability that  $\bar{X}(i) = 0$  can then be estimated from above by the expected number of Toom contours rooted at  $i$  that are present in  $\Phi$ . In particular, applying this to the maximal trajectory  $\bar{X}^p$  of the random monotone cellular automaton  $\Phi^p$ , we are under certain additional assumptions able to show that the density  $\bar{\rho}(p) = \mathbb{P}[\bar{X}^p(i, t) = 1]$  of the upper invariant law, which in this case does not depend on  $(i, t)$ , tends to one as  $p \rightarrow 0$ . In its essence, the method goes back to Toom's proof of [Too80, Thm 5] but we have significantly modified and simplified the argument with the aim of making it more flexible and intuitive. One of the most significant changes we have made is the introduction of sources and sinks (see Figure 2 below). By contrast, the contours used in [Too80] are directed graphs in which the number of incoming edges equals the number of outgoing edges at each vertex.

## 2.2 Toom contours

We will need directed graphs in which both the vertices and the edges can have different types. Let  $A$  and  $B$  be finite sets. By definition, a *typed directed graph* with *vertex set*  $V$ , *vertex type set*  $A$ , and *edge type set*  $B$  is a pair  $(\mathcal{V}, \mathcal{E})$  where  $\mathcal{V}$  is a subset of  $V \times A$  and  $\mathcal{E}$  is a subset of  $V \times V \times B$ , such that

$$\forall v \in V \exists a \in A \text{ s.t. } (v, a) \in \mathcal{V}. \quad (2.6)$$

For each  $a \in A$  and  $b \in B$ , we call

$$V_a := \{v : (v, a) \in \mathcal{V}\} \quad \text{and} \quad \vec{E}_b := \{(v, w) : (v, w, b) \in \mathcal{E}\} \quad (2.7)$$

the set of vertices of type  $a$  and the set of directed edges of type  $b$ , respectively. Note that vertices can have more than one type, i.e.,  $V_a$  and  $V_{a'}$  are not necessarily disjoint for  $a \neq a'$ , and the same applies to edges. As a consequence, several edges of different types can connect the same two vertices  $v, w$ , but always at most one of each type. If  $(\mathcal{V}, \mathcal{E})$  is a typed directed graph, then we let  $(V, \vec{E})$  denote the directed graph given by

$$V = \bigcup_{a \in A} V_a \quad \text{and} \quad \vec{E} := \bigcup_{b \in B} \vec{E}_b, \quad (2.8)$$

where the first equality follows from (2.6) and the second equality is a definition. We call  $(V, \vec{E})$  the *untyped* directed graph associated with  $(\mathcal{V}, \mathcal{E})$ . We also set  $E := \{(v, w) : (v, w) \in \vec{E}\}$ . Then  $(V, E)$  is an undirected graph, which we call the undirected graph *associated with*  $(V, \vec{E})$ . We say that a typed directed graph  $(\mathcal{V}, \mathcal{E})$  or a directed graph  $(V, \vec{E})$  are *connected* if their associated undirected graph  $(V, E)$  is connected. A *rooted* directed graph is a triple  $(v_\circ, V, \vec{E})$  such that  $(V, \vec{E})$  is a directed graph and  $v_\circ \in V$  is a specially designated vertex, called the *root*. Rooted undirected graphs and rooted typed directed graphs are defined in the same way.

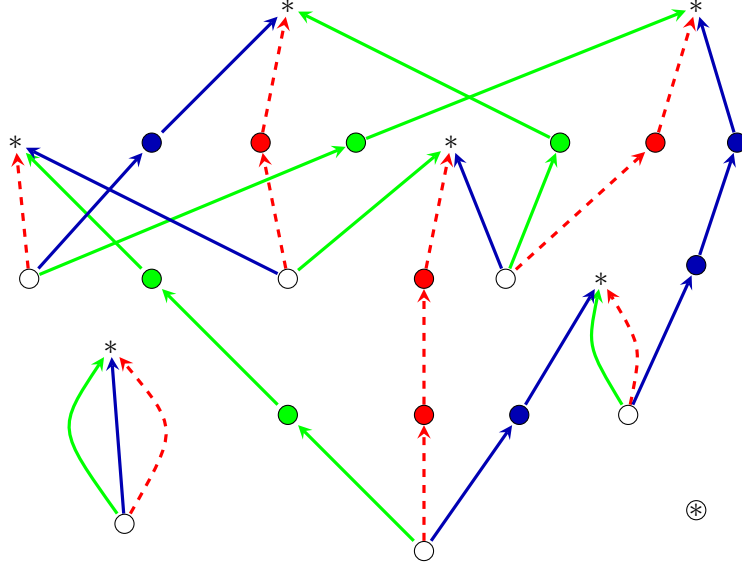


Figure 2: Example of a Toom graph with three charges. Sources are indicated with open dots, sinks with asterisks, and internal vertices and edges of the three possible charges with three colours. Note the isolated vertex in the lower right corner, which is a source and a sink at the same time.

For any directed graph  $(V, \vec{E})$ , we let

$$\vec{E}_{\text{in}}(v) := \{(u, v') \in \vec{E} : v' = v\} \quad \text{and} \quad \vec{E}_{\text{out}}(v) := \{(v', w) \in \vec{E} : v' = v\} \quad (2.9)$$

denote the sets of directed edges entering and leaving a given vertex  $v \in V$ , respectively. Similarly, in a typed directed graph,  $\vec{E}_{b,\text{in}}(v)$  and  $\vec{E}_{b,\text{out}}(v)$  denote the sets of incoming or outgoing directed edges of type  $b$  at  $v$ .

We adopt the following general notation. For any directed graph  $(V, \vec{E})$ , set  $\Lambda$ , and function  $\psi : V \rightarrow \Lambda$ , we let

$$\psi(V) := \{\psi(v) : v \in V\} \quad \text{and} \quad \psi(\vec{E}) := \{(\psi(v), \psi(w)) : (v, w) \in \vec{E}\} \quad (2.10)$$

denote the images of  $V$  and  $\vec{E}$  under  $\psi$ . We can naturally view  $(\psi(V), \psi(\vec{E}))$  as a directed graph with set of vertices  $\psi(V)$  and set of directed edges  $\psi(\vec{E})$ . We denote this graph by  $\psi(V, \vec{E}) := (\psi(V), \psi(\vec{E}))$ . Similarly, if  $(\mathcal{V}, \mathcal{E})$  is a typed directed graph, then we let  $\psi(\mathcal{V}, \mathcal{E})$  denote the typed directed graph defined as

$$\begin{aligned} \psi(\mathcal{V}, \mathcal{E}) &:= (\psi(\mathcal{V}), \psi(\mathcal{E})) & \text{with} & \quad \psi(\mathcal{V}) := \{(\psi(v), a) : (v, a) \in \mathcal{V}\} \\ & & \text{and} & \quad \psi(\mathcal{E}) := \{(\psi(v), \psi(w), b) : (v, w, b) \in \mathcal{E}\}. \end{aligned} \quad (2.11)$$

Also, if  $(v_\circ, V, \vec{E})$  is a rooted directed graph, then we let  $\psi(v_\circ, V, \vec{E})$  denote the rooted directed graph  $(\psi(v_\circ), \psi(V), \psi(\vec{E}))$ , and we use similar notation for rooted typed directed graphs. Two typed directed graphs  $(\mathcal{V}, \mathcal{E})$  and  $(\mathcal{W}, \mathcal{F})$  are *isomorphic* if there exists a bijection  $\psi : V \rightarrow W$  such that  $\psi(\mathcal{V}, \mathcal{E}) = (\mathcal{W}, \mathcal{F})$ . Similar conventions apply to directed graphs, rooted directed graphs, and so on.

**Definition 16** A Toom graph with  $\sigma \geq 1$  charges is a typed directed graph  $(\mathcal{V}, \mathcal{E})$  with vertex type set  $\{\circ, *, 1, \dots, \sigma\}$  and edge type set  $\{1, \dots, \sigma\}$  that satisfies the following conditions:

- (i)  $|\vec{E}_{s,\text{in}}(v)| = 0$  ( $1 \leq s \leq \sigma$ ) and  $|\vec{E}_{1,\text{out}}(v)| = \dots = |\vec{E}_{\sigma,\text{out}}(v)| \leq 1$  for all  $v \in V_\circ$ .



- (ii)  $|\vec{E}_{s,\text{out}}(v)| = 0$  ( $1 \leq s \leq \sigma$ ) and  $|\vec{E}_{1,\text{in}}(v)| = \dots = |\vec{E}_{\sigma,\text{in}}(v)| \leq 1$  for all  $v \in V_*$ .
- (iii)  $|\vec{E}_{s,\text{in}}(v)| = 1 = |\vec{E}_{s,\text{out}}(v)|$  and  $|\vec{E}_{l,\text{in}}(v)| = 0 = |\vec{E}_{l,\text{out}}(v)|$  for each  $l \neq s$  and  $v \in V_s$ .

See Figure 2 for a picture of a Toom graph with three charges. Vertices in  $V_\circ$ ,  $V_*$ , and  $V_s$  are called *sources*, *sinks*, and *internal vertices* with *charge*  $s$ , respectively. Vertices in  $V_\circ \cap V_*$  are called *isolated vertices*. With the exception of isolated vertices, the inequalities  $\leq 1$  in (i) and (ii) are equalities. Informally, we can imagine that at each source there emerge  $\sigma$  charges, one of each type, that then travel via internal vertices of the corresponding charge through the graph until they arrive at a sink, in such a way that at each sink there converge precisely  $\sigma$  charges, one of each type. This informal picture holds even for isolated vertices, if we imagine that in this case, the charges arrive immediately at the sink that is at the same time a source. It is clear from this informal picture that  $|V_\circ| = |V_*|$ , i.e., the number of sources equals the number of sinks. We let  $(V, \vec{E})$  denote the directed graph associated with  $(\mathcal{V}, \mathcal{E})$ .

Toom graphs and the Toom contours that will be defined below were designed to make use of the condition (1.21) of Theorem 9 on the worst-case edge speeds. The curious reader may skip ahead to the beginning of Subsection 5.3 where we give an informal description of the main idea of the proof of Theorem 9.

Recall that a rooted directed graph is a directed graph with a specially designated vertex, called the root. In the case of Toom graphs, we will always assume that the root is a source.

**Definition 17** *A rooted Toom graph with  $\sigma \geq 1$  charges is a rooted typed directed graph  $(v_\circ, \mathcal{V}, \mathcal{E})$  such that  $(\mathcal{V}, \mathcal{E})$  is a Toom graph with  $\sigma \geq 1$  charges and  $v_\circ \in V_\circ$ . For any rooted Toom graph  $(v_\circ, \mathcal{V}, \mathcal{E})$ , we write*

$$V'_\circ := V_\circ \setminus \{v_\circ\} \quad \text{and} \quad V'_s := V_s \cup \{v_\circ\} \quad (1 \leq s \leq \sigma). \quad (2.12)$$

The idea behind (2.12) is that for rooted Toom contours, we view the root more as if it were a collection of internal vertices than as a source. This is reflected in condition (ii) of the following definition.

**Definition 18** *Let  $(v_\circ, \mathcal{V}, \mathcal{E})$  be a rooted Toom graph and let  $\Lambda$  be a countable set. An embedding of  $(v_\circ, \mathcal{V}, \mathcal{E})$  in  $\Lambda$  is a map  $\psi : V \rightarrow \Lambda$  such that:*

- (i)  $\psi(v_1) \neq \psi(v_2)$  for each  $v_1 \in V_*$  and  $v_2 \in V$  with  $v_1 \neq v_2$ ,
- (ii)  $\psi(v_1) \neq \psi(v_2)$  for each  $v_1, v_2 \in V'_s$  with  $v_1 \neq v_2$  ( $1 \leq s \leq \sigma$ ).

Condition (i) says that sinks do not overlap with other vertices and condition (ii) says that internal vertices do not overlap with other internal vertices of the same charge, where in line with (2.12) we view the root as a collection of internal vertices. We make the following observation.

**Lemma 19 (No double incoming edges)** *Let  $\psi$  be an embedding of a rooted Toom graph  $(v_\circ, \mathcal{V}, \mathcal{E})$  with  $\sigma \geq 1$  edges in a set  $\Lambda$ . Then*

$$|\{(v, w) \in \vec{E}_s : \psi(w) = j\}| \leq 1 \quad (j \in \Lambda, 1 \leq s \leq \sigma). \quad (2.13)$$

**Proof** Immediate from Definition 18, since each charged edge ends in an internal vertex of the same charge or in a sink. ■

**Definition 20** *Let  $\Lambda$  be a countable set. A Toom contour in  $\Lambda$  with  $\sigma \geq 1$  charges is a quadruple  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi)$ , where  $(v_\circ, \mathcal{V}, \mathcal{E})$  is a rooted connected Toom graph with  $\sigma$  charges and  $\psi$  is an embedding of  $(v_\circ, \mathcal{V}, \mathcal{E})$  in  $\Lambda$ . We say that the Toom contour is rooted at  $i_\circ := \psi(v_\circ)$ . Two Toom contours  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  and  $(v'_\circ, \mathcal{V}', \mathcal{E}', \psi')$  are isomorphic if there exists a bijection  $\chi : V \rightarrow V'$  such that  $\chi(v_\circ, \mathcal{V}, \mathcal{E}) = (v'_\circ, \mathcal{V}', \mathcal{E}')$  and  $\psi(v) = \psi'(\chi(v))$  ( $v \in V$ ). We say that  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  and  $(v'_\circ, \mathcal{V}', \mathcal{E}', \psi')$  are equivalent if, using notation introduced in (2.11), one has  $\psi(\mathcal{V}, \mathcal{E}) = \psi'(\mathcal{V}', \mathcal{E}')$ .*

We note that as a result of Lemma 19, each charged edge in  $\psi(\mathcal{E})$  corresponds to a unique charged edge in  $\mathcal{E}$ . Two isomorphic Toom contours are clearly equivalent, but the converse implication does not hold, since sources can overlap with each other and with internal vertices and as a result, although two equivalent Toom contours have charged edges in the same locations, these edges can be differently connected leading to two Toom contours that are not isomorphic. See Figure 3 for an example of a Toom contour with two charges.

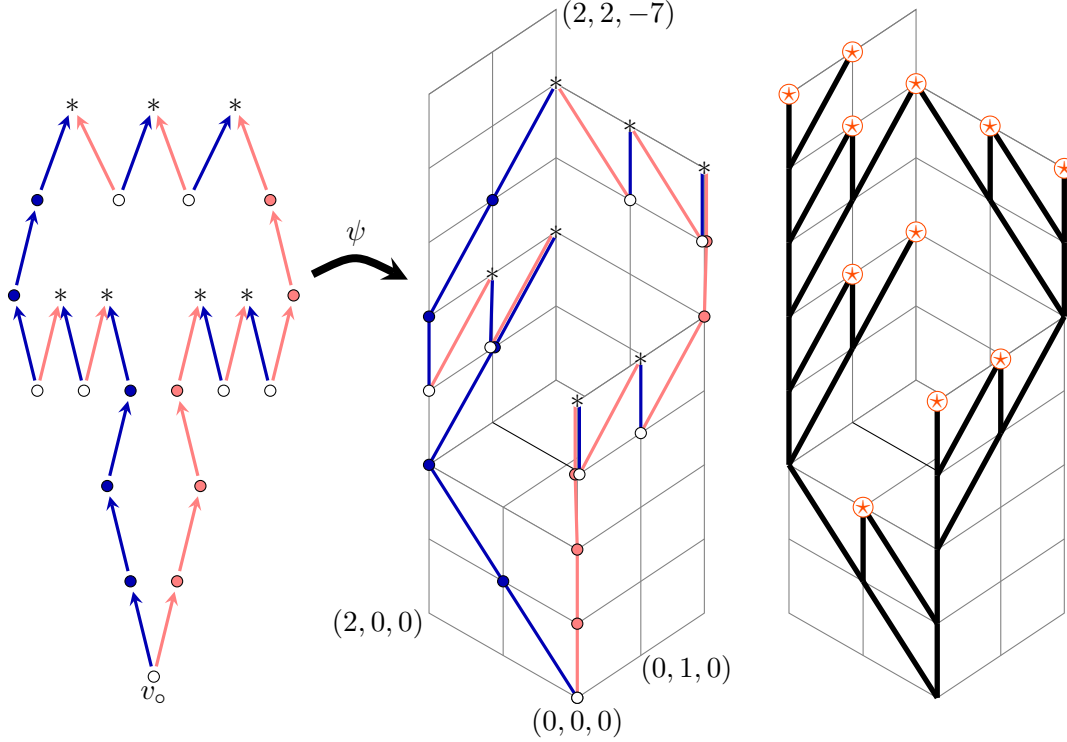


Figure 3: A Toom contour in  $\mathbb{Z}^3$  rooted at  $(0, 0, 0)$ . The third coordinate represents time and is plotted downwards. The picture on the right shows a minimal explanation (or rather its associated undirected explanation graph as defined in Subsection 4.1) for a monotone cellular automaton  $\Phi^p$  that applies the maps  $\varphi^0$  and  $\varphi^{\text{coop}}$  with probabilities  $p$  and  $1 - p$ , respectively. The origin has the value zero because the sites marked with a star are defective; removing any of these defective sites results in the origin having the value one. The Toom contour in the middle picture is present in  $\Phi^p$ . In particular, the sinks of the Toom contour coincide with some, though not with all of the defective sites of the minimal explanation.

### 2.3 Presence of Toom contours

Our next aim is to define when a Toom contour is present in a monotone cellular automaton  $\phi = (\phi_i)_{i \in \Lambda}$ . This will require us to make some extra assumptions and equip  $\phi$  with some extra structure.

**Definition 21** A typed dependence graph with  $\sigma \geq 1$  types of edges is a typed directed graph  $(\Lambda, \mathcal{H})$  with vertex type set  $\{0, 1, \bullet\}$  and edge type set  $\{1, \dots, \sigma\}$  such that for  $\vec{H}_s := \{(i, j) : (i, j, s) \in \mathcal{H}\}$

- (i)  $\vec{H}_{s, \text{out}}(i) = \emptyset$  for all  $i \in \Lambda_0 \cup \Lambda_1$  and  $1 \leq s \leq \sigma$ ,
- (ii)  $\vec{H}_{s, \text{out}}(i) \neq \emptyset$  for all  $i \in \Lambda_\bullet$  and  $1 \leq s \leq \sigma$ ,

and its associated untyped directed graph  $(\Lambda, \vec{H})$  is acyclic. The monotone cellular automaton  $\phi = (\phi_i)_{i \in \Lambda}$  associated with the typed dependence graph  $(\Lambda, \mathcal{H})$  is defined by

$$\phi_i(x) = \begin{cases} \bigvee_{s=1}^{\sigma} \bigwedge_{j: (i,j) \in \vec{H}_s} x(j) & \text{if } i \in \Lambda_{\bullet}, \\ r & \text{if } i \in \Lambda_r \quad (r = 0, 1), \end{cases} \quad (2.14)$$

$(i \in \Lambda, x \in \{0, 1\}^{\Lambda})$ .

It is easy to see that if  $(\Lambda, \mathcal{H})$  is a typed dependence graph,  $\phi$  is its associated monotone cellular automaton, and  $(\Lambda, \vec{H})$  is its associated untyped directed graph, then  $(\Lambda, \vec{H})$  is the dependence graph of  $\phi$  as defined in Subsection 2.1. In particular, the assumption that  $(\Lambda, \vec{H})$  is acyclic guarantees that (2.14) defines a cellular automaton. It is clear from (2.14) that  $\phi_i$  is monotone for each  $i \in \Lambda$  and that  $\phi_i$  is one of the constant maps  $\phi^r$  ( $r = 0, 1$ ) defined in (2.1) if and only if  $i \in \Lambda_r$  ( $r = 0, 1$ ). Furthermore, recalling (1.9) and the definition of one-sets, we can see that for each  $i \in \Lambda_{\bullet}$  the sets  $\{j \in \Lambda : (i, j) \in \vec{H}_s\}$  ( $1 \leq s \leq \sigma$ ) are one-sets of  $\phi_i$ , though not necessarily minimal. Elements of  $\Lambda_0$ , where the constant zero map is applied, are called *defective sites*. Below, we make use of the definition (2.12), i.e., we treat the root as if it were a collection of internal vertices.

**Definition 22** *Let  $(\Lambda, \mathcal{H})$  be a typed dependence graph with  $\sigma \geq 1$  types of edges. We say that a Toom contour  $(v_{\circ}, \mathcal{V}, \mathcal{E}, \psi)$  with  $\sigma$  charges is present in  $(\Lambda, \mathcal{H})$  if:*

- (i)  $\psi(v) \in \Lambda_0$  for all  $v \in V_*$ ,
- (ii)  $(\psi(v), \psi(w)) \in \vec{H}_s$  for all  $(v, w) \in \vec{E}_s^{\bullet}$  ( $1 \leq s \leq \sigma$ ),
- (iii)  $(\psi(v), \psi(w)) \in \vec{H}$  for all  $(v, w) \in \vec{E}^{\circ}$ ,

where for any rooted Toom graph  $(v_{\circ}, \mathcal{V}, \mathcal{E})$ , we write

$$\begin{aligned} \vec{E}^{\bullet} &:= \bigcup_{s=1}^{\sigma} \vec{E}_s^{\bullet} \quad \text{with} \quad \vec{E}_s^{\bullet} := \{(v, w) \in \vec{E}_s : v \in V'_s\} \quad (1 \leq s \leq \sigma), \\ \vec{E}^{\circ} &:= \bigcup_{s=1}^{\sigma} \vec{E}_s^{\circ} \quad \text{with} \quad \vec{E}_s^{\circ} := \{(v, w) \in \vec{E}_s : v \in V'_s\} \quad (1 \leq s \leq \sigma). \end{aligned} \quad (2.15)$$

Condition (i) says that sinks of the Toom contour correspond to defective sites of the typed dependence graph. Conditions (ii) and (iii) say that directed edges of the Toom graph  $(\mathcal{V}, \mathcal{E})$  are mapped to directed edges of the typed dependence graph  $(\Lambda, \mathcal{H})$ , where edges coming out of an internal vertex must be mapped to edges of the corresponding type, and we treat the root as if it were a collection of internal vertices. Let  $W := \psi(V)$  and  $W_* := \psi(V_*)$ . We note that Definition 22 implies that

$$W \cap \Lambda_0 = W_* \quad \text{and} \quad W \cap \Lambda_1 = \emptyset. \quad (2.16)$$

Indeed, the inclusion  $W_* \subset W \cap \Lambda_0$  is immediate from condition (i) while conditions (ii) and (iii) imply  $W \setminus W_* \subset \Lambda \setminus (\Lambda_0 \cup \Lambda_1)$  since for  $v \in V \setminus V_*$ , one has  $\vec{E}_{\text{out}}(v) \neq \emptyset$ , while  $\vec{H}_{\text{out}}(i) = \emptyset$  for  $i \in \Lambda_0 \cup \Lambda_1$ .

The following crucial theorem, proved in Subsection 4.2, links the maximal trajectory to Toom contours. In its essence, this goes back to part 3 of the proof of [Too80, Thm 1], but we have reformulated things to a point where, putting the two texts besides each other, it is hard at first sight to spot the similarity.

**Theorem 23 (Presence of a Toom contour)** *Let  $(\Lambda, \mathcal{H})$  be a typed dependence graph with  $\sigma \geq 1$  types of edges, let  $\phi$  be its associated monotone cellular automaton, and let  $\bar{x}$  be its maximal trajectory. If  $\bar{x}(i) = 0$  for some  $i \in \Lambda$ , then a Toom contour  $(v_{\circ}, \mathcal{V}, \mathcal{E}, \psi)$  rooted at  $i$  is present in  $(\Lambda, \mathcal{H})$ .*

We note that the converse of Theorem 23 does not hold, i.e., the presence in  $(\Lambda, \mathcal{H})$  of a Toom contour  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  does not imply that  $\bar{x}(i) = 0$ . This can be seen from Figure 3. In this example, if there would be no other defective sites apart from the sinks of the Toom contour, then the origin would have the value one. This is a difference with the Peierls arguments used in percolation theory, where the presence of a contour is a necessary and sufficient condition for the absence of percolation.

## 2.4 Toom cycles

We will give two proofs of Theorem 23: one that works for general  $\sigma \geq 1$ , and another that works only for  $\sigma = 2$ , but that in this case gives some extra information that can sometimes be used to get sharper bounds. As Figure 3 shows, Toom contours with two charges are essentially cycles. In the present subsection, we define Toom cycles, which are Toom contours with two charges that have some useful additional properties, and we formulate a theorem about the presence of Toom cycles in monotone cellular automata.

For  $n \geq 2$ , let  $[n] := \{0, \dots, n-1\}$  equipped with addition modulo  $n$ . We define a *cycle* of length  $n \geq 2$  to be an undirected graph  $(V, E)$  with vertex set  $V = [n]$  and edge set  $E := \{\{v, v+1\} : v \in [n]\}$ . Similarly, we define a cycle of length 1 to be the undirected graph  $(V, E) := (\{0\}, \emptyset)$ . We define an *oriented cycle* of length  $n \geq 1$  to be a directed graph  $(V, \vec{E})$  whose associated undirected graph  $(V, E)$  is a cycle of length  $n$  such that for each undirected edge  $\{v, w\} \in E$ , precisely one of the directed edges  $(v, w)$  and  $(w, v)$  is an element of  $\vec{E}$  (but not both). In other words, this is a cycle in which each undirected edge has been given an orientation.

Each oriented cycle of length  $n \geq 2$  naturally gives rise to a connected Toom graph  $(\mathcal{V}, \mathcal{E})$  with two charges by setting

$$\begin{aligned} V_\circ &:= \{v \in [n] : (v, v-1), (v, v+1) \in \vec{E}\}, \\ V_* &:= \{v \in [n] : (v-1, v), (v+1, v) \in \vec{E}\}, \\ V_1 &:= \{v \in [n] : (v-1, v), (v, v+1) \in \vec{E}\}, \\ V_2 &:= \{v \in [n] : (v+1, v), (v, v-1) \in \vec{E}\} \end{aligned} \tag{2.17}$$

and

$$\vec{E}_1 := \{(v, w) \in \vec{E} : w = v+1\} \quad \text{and} \quad \vec{E}_2 := \{(v, w) \in \vec{E} : w = v-1\}. \tag{2.18}$$

Similarly, we may associate the oriented cycle of length one with the trivial Toom graph  $(\mathcal{V}, \mathcal{E})$  with two charges defined as  $V_\circ = V_* := \{0\}$  and  $V_1 = V_2 = \vec{E}_1 = \vec{E}_2 := \emptyset$ . If  $0 \in V_\circ$ , then we can take  $v_\circ := 0$  to be the root. In view of this, connected rooted Toom graphs with two charges correspond (up to isomorphism) precisely to oriented cycles  $(V, \vec{E})$  of length  $n \geq 1$  for which  $0 \in V_\circ$ .

It is sometimes convenient to add the element  $n$  to  $V$  and to replace the oriented edge  $(0, n-1) \in \vec{E}_2$  by  $(n, n-1)$ . Thus, we may identify an oriented cycle  $(V, \vec{E})$  of length  $n \geq 1$  for which  $0 \in V_\circ$  with an oriented path of length  $n$  for which  $(0, 1) \in \vec{E}_1$  and  $(n, n-1) \in \vec{E}_2$ . Similar to what we did in (2.12), it will be convenient to define

$$V'_\circ := V_\circ \setminus \{0\}, \quad V'_1 := V_1 \cup \{0\} \quad \text{and} \quad V'_2 := V_2 \cup \{n\}. \tag{2.19}$$

In condition (ii) of the following definition, we equip  $\{0, \dots, n\}$  with the natural total order and we equip  $\{1, \circ, 2\}$  with the total order  $1 < \circ < 2$ .

**Definition 24** *Let  $\Lambda$  be a countable set. A Toom cycle in  $\Lambda$  is a triple  $(V, \vec{E}, \psi)$  where  $(V, \vec{E})$  is an oriented cycle of length  $n \geq 1$  such that  $0 \in V_\circ$ , and  $\psi : \{0, \dots, n\} \rightarrow \Lambda^6$  is a map such that  $\psi_0 = \psi_n$  and*

<sup>6</sup>For ease of notation we write  $\psi_v$  for  $\psi(v)$  in case of Toom cycles.

- (i)  $\psi_v \neq \psi_w$  for each  $v \in V_*$  and  $w \in V$  with  $v \neq w$ ,
- (ii) if  $\psi_v = \psi_w$  for some  $v \in V'_s$  and  $w \in V'_t$  with  $s, t \in \{1, \circ, 2\}$  and  $s \leq t$ , then  $v \leq w$ .

We say that the Toom cycle  $(V, \vec{E}, \psi)$  is rooted at  $\psi_0$ .

If  $(V, \vec{E}, \psi)$  is a Toom cycle and  $(v_\circ, \mathcal{V}, \mathcal{E})$  is its associated rooted Toom graph, then  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  is a Toom contour with two charges. We call this the Toom contour *associated with* the Toom cycle  $(V, \vec{E}, \psi)$ . Applying property (ii) with  $s = t$  implies

- if  $\psi_v = \psi_w$  for some  $v, w \in V'_s$  with  $s \in \{1, \circ, 2\}$ , then  $v = w$ ,

so property (ii) of Definition 24 implies property (ii) of Definition 18. It is easy to see that it is strictly stronger, so not every Toom contour with two charges comes from a Toom cycle.

We next define what it means for a Toom cycle to be present in a typed dependence graph  $(\Lambda, \mathcal{H})$ .

**Definition 25** *Let  $(\Lambda, \mathcal{H})$  be a typed dependence graph with 2 types of edges as in Definition 21. We say that a Toom cycle  $(V, \vec{E}, \psi)$  of length  $n \geq 2$  is present in  $(\Lambda, \mathcal{H})$  if:*

- (i)  $\psi_v \in \Lambda_0$  for all  $v \in V_*$ ,
- (ii)  $(\psi_v, \psi_w) \in \vec{H}_s$  for all  $(v, w) \in \vec{E}_s$  with  $v \in V'_s$  ( $s = 1, 2$ ),
- (iii)  $(\psi_v, \psi_w) \in \vec{H}_{3-s}$  for all  $(v, w) \in \vec{E}_s$  with  $v \in V'_\circ$  ( $s = 1, 2$ ).

For Toom cycles of length 1, only condition (i) applies.

Note that in condition (iii) above,  $3 - s$  is 1 if  $s = 2$  and 2 if  $s = 1$ , so this condition says that directed edges coming out of a source other than the root must use a directed edge of  $(\Lambda, \mathcal{H})$  of the opposite charge. This condition is stronger than condition (iii) of Definition 22. One can check that our definition implies that if a Toom cycle is present in  $(\Lambda, \mathcal{H})$ , then its associated Toom contour is present in  $(\Lambda, \mathcal{H})$  in the sense of Definition 22, but because of our previous remark, the converse implication does not hold. One can check that the Toom contour with two charges in Figures 3 and 4 comes from a Toom cycle that is present in the strong sense of Definition 25.

In the same way as in (2.16), one can see that Definition 25 implies that  $\psi_v \notin \Lambda_1$  for all  $v \in V$ . In view of our previous remarks, the following theorem strengthens Theorem 23 in the special case of two charges. Our proof of Theorem 26 (in Subsection 4.3) will largely be independent of the proof of Theorem 23.

**Theorem 26 (Presence of a Toom cycle)** *Let  $(\Lambda, \mathcal{H})$  be a typed dependence graph with 2 types of edges, let  $\phi$  be its associated monotone cellular automaton, and let  $\bar{x}$  be its maximal trajectory. If  $\bar{x}(i) = 0$  for some  $i \in \Lambda$ , then a Toom cycle rooted at  $i$  is present in  $(\Lambda, \mathcal{H})$ .*

## 2.5 A Peierls bound

Theorems 23 and 26 can be used to prove upper bounds on the probability that the maximal trajectory of a random monotone cellular automaton takes the value zero in a given point. For concreteness, we formulate this as a theorem.

**Theorem 27 (Peierls bound)** *Let  $\Phi = (\Phi_i)_{i \in \Lambda}$  be a random monotone cellular automaton and let  $\bar{X}$  be its maximal trajectory. Let*

$$\Lambda_r := \{i \in \Lambda : \Phi_i = \phi^r\} \quad (r = 0, 1) \quad \text{and} \quad \Lambda_\bullet := \Lambda \setminus (\Lambda_0 \cup \Lambda_1). \quad (2.20)$$

Let  $\sigma \geq 1$  be an integer and for each  $i \in \Lambda_\bullet$  and  $1 \leq s \leq \sigma$ , let  $A_{s,i} \in \mathcal{A}(\Phi_i)$ . Define  $\mathcal{H} = (\vec{H}_1, \dots, \vec{H}_\sigma)$  by

$$\vec{H}_s := \{(i, j) : i \in \Lambda_\bullet, j \in A_{s,i}\} \quad (1 \leq s \leq \sigma). \quad (2.21)$$

Fix  $i \in \Lambda$  and let  $\mathcal{T}_i$  denote the set of Toom contours rooted at  $i$  (up to equivalence). Then

$$\mathbb{P}[\overline{X}(i) = 0] \leq \sum_{T \in \mathcal{T}_i} \mathbb{P}[T \text{ is present in } (\Lambda, \mathcal{H})]. \quad (2.22)$$

If  $\sigma = 2$ , then (2.22) remains true if we restrict the sum to Toom cycles rooted at  $i$ .

**Proof** Let  $\Psi$  be the random monotone cellular automaton associated with the random typed dependence graph  $(\Lambda, \mathcal{H})$ . Then in view of (2.14) we have  $\Psi_i = \Phi_i$  for  $i \in \Lambda_0 \cup \Lambda_1$  while

$$\Psi_i(x) = \bigvee_{s=1}^{\sigma} \bigwedge_{j \in A_{s,i}} x(j) \leq \bigvee_{A \in \mathcal{A}(\Phi_i)} \bigwedge_{j \in A} x(j) = \Phi_i(x) \quad (i \in \Lambda_\bullet, x \in \{0, 1\}^\Lambda). \quad (2.23)$$

Using this, it is easy to check (see Lemma 28 below) that the maximal trajectories  $\overline{X}$  of  $\Phi$  and  $\overline{Y}$  of  $\Psi$  are ordered as  $\overline{Y} \leq \overline{X}$  (pointwise). In particular,

$$\mathbb{P}[\overline{X}(i) = 0] \leq \mathbb{P}[\overline{Y}(i) = 0]. \quad (2.24)$$

By Theorems 23 and 26, the right-hand side of (2.24) can be bounded from above by the probability that there is a Toom contour or cycle present in  $(\Lambda, \mathcal{H})$ , which in turn can be estimated from above by the expected number of Toom contours or cycles.  $\blacksquare$

## Part II

# Proofs

### 3 Preliminaries

#### 3.1 Eroders

In this subsection we prove Lemmas 7 and 10.

**Proof of Lemma 7** It suffices to prove the claim for  $t = 1$ . Fix  $j \in \mathbb{Z}^d$  and set  $j + A := \{j + i : i \in A\}$  ( $A \in \mathcal{A}(\varphi)$ ). Then one has  $\Psi_{0,t}(H_r^\ell)(j) = 1$  if and only if there exists an  $A \in \mathcal{A}(\varphi)$  such that  $\ell(k) \geq r$  for all  $k \in j + A$ . Equivalently, this says that

$$\sup_{A \in \mathcal{A}(\varphi)} \inf_{k \in j + A} \ell(k) \geq r. \quad (3.1)$$

Using (1.15) and linearity, we can rewrite this as  $\ell(j) + \varepsilon_\varphi(\ell) \geq r$ , which is equivalent to  $j \in H_{r - \varepsilon_\varphi(\ell)}^\ell$ .  $\blacksquare$

We next prove Lemma 10. Our proof depends on the equivalence of (1.10) and the eroder property, which is proved in [Pon13, Thm 1]. We recall that the fact that (1.10) implies the eroder property has already been demonstrated below Lemma 10, so we depend on [Pon13, Thm 1] only for the converse implication.

**Proof of Lemma 10** In [Pon13, Lemma 12] it is shown<sup>7</sup> that (1.10) is equivalent to the existence of a linear polar function  $L$  of dimension  $2 \leq \sigma \leq d + 1$  and constants  $\varepsilon_1, \dots, \varepsilon_\sigma$  such

<sup>7</sup>Since Ponselet discusses stability of the all-zero fixed point while we discuss stability of the all-one fixed point, in [Pon13] the roles of zeros and ones are reversed compared to our conventions.

that  $\sum_{s=1}^{\sigma} \varepsilon_s > 0$  and for each  $1 \leq s \leq \sigma$ , there exists an  $A_s \in \mathcal{A}(\varphi)$  such that  $\varepsilon_s - L_s(i) \leq 0$  for all  $i \in A_s$ . It follows that

$$\sum_{s=1}^{\sigma} \sup_{A \in \mathcal{A}(\varphi)} \inf_{i \in A} L_s(i) \geq \sum_{s=1}^{\sigma} \inf_{i \in A_s} L_s(i) \geq \sum_{s=1}^{\sigma} \varepsilon_s > 0, \quad (3.2)$$

which shows that (1.22) holds. Assume, conversely, that (1.22) holds. Since  $\mathcal{A}(\varphi)$  is finite, for each  $1 \leq s \leq \sigma$  we can choose  $A_s \in \mathcal{A}(\varphi)$  such that

$$\varepsilon_s := \inf_{i \in A_s} L_s(i) = \sup_{A \in \mathcal{A}(\varphi)} \inf_{i \in A} L_s(i). \quad (3.3)$$

Then (1.22) says that  $\sum_{s=1}^{\sigma} \varepsilon_s > 0$ . Let  $H_s := \{z \in \mathbb{R}^d : L_s(z) \geq \varepsilon_s\}$ . By the definition of a linear polar function,  $\sum_{s=1}^{\sigma} L_s(z) = 0$  for each  $z \in \mathbb{R}^d$ , and hence the condition  $\sum_{s=1}^{\sigma} \varepsilon_s > 0$  implies that for each  $z \in \mathbb{R}^d$ , there exists an  $1 \leq s \leq \sigma$  such that  $L_s(z) < \varepsilon_s$ . In other words, this says that  $\bigcap_{s=1}^{\sigma} H_s = \emptyset$ . For each  $1 \leq s \leq \sigma$ , the set  $A_s$  is contained in the half-space  $H_s$  and hence the same is true for  $\text{Conv}(A_s)$ , so we conclude that

$$\bigcap_{s=1}^{\sigma} \text{Conv}(A_s) = \emptyset, \quad (3.4)$$

from which (1.10) follows. ■

### 3.2 The maximal trajectory

In this subsections, we prove Lemmas 14 and 15, as well as Lemma 28 that has already been used in the proof of Theorem 27.

**Proof of Lemma 14** By symmetry, it suffices to show that there exists a trajectory  $\bar{x}$  that is uniquely characterised by the property that each trajectory  $x$  of  $\phi$  satisfies  $x \leq \bar{x}$ . Let  $\Lambda_n \subset \Lambda$  be finite sets increasing to  $\Lambda$  and for each  $n$ , let  $\phi^n$  denote the monotone cellular automaton defined by

$$\phi_i^n := \begin{cases} \phi^1 & \text{if } i \in \Lambda \setminus \Lambda_n \\ \phi_i & \text{if } i \in \Lambda_n, \end{cases} \quad (3.5)$$

where  $\phi^1$ , defined in (2.1), denotes the map that is constantly one. Since  $\Lambda_n$  is finite, it is easy to see that  $\phi^n$  has a unique trajectory  $x^n$ , which satisfies  $x^n(i) = 1$  for all  $i \in \Lambda \setminus \Lambda_n$ . One has  $x^n \geq x^{n+1}$  (coordinatewise) for each  $n$  so the monotone limit  $\bar{x}(i) := \lim_{n \rightarrow \infty} x^n(i)$  ( $i \in \Lambda$ ) exists. It is straightforward to check that  $\bar{x}$  is a trajectory of  $\phi$ . If  $x$  is any other trajectory of  $\phi$ , then  $x \leq x^n$  for all  $n$  and hence  $x \leq \bar{x}$ . ■

**Proof of Lemma 15** By symmetry, it suffices to prove the claim for the upper invariant law. For each  $n \geq 0$ , let  $\Phi^{n,p}$  denote the modified cellular automaton defined by

$$\Phi_{i,t}^{n,p} := \begin{cases} \varphi^1 & \text{if } t \leq -n \\ \Phi_{i,t}^p & \text{if } t > -n. \end{cases} \quad (3.6)$$

Then it is easy to see that  $\Phi^{n,p}$  has a unique trajectory  $X^{n,p}$ , which satisfies  $X^{n,p}(i, t) = 1$  for all  $i \in \mathbb{Z}^d$  and  $t \leq -n$ . Exactly the same argument as in the proof of Lemma 14 shows that  $X^{n,p} \rightarrow \bar{X}^p$  (pointwise) almost surely. The claim now follows from the observation that  $(X^{n,p}(i, 0))_{i \in \mathbb{Z}^d}$  is equally distributed with the random variable  $X_n^p$  in the second formula of (1.4). ■

**Lemma 28 (Comparison of maximal trajectories)** *Let  $\phi = (\phi_i)_{i \in \Lambda}$  and  $\psi = (\psi_i)_{i \in \Lambda}$  be monotone cellular automata and let  $\bar{x}$  and  $\bar{y}$  denote their respective maximal trajectories. Assume that*

$$\phi_i(x) \leq \psi_i(x) \quad (i \in \Lambda, x \in \{0, 1\}^\Lambda). \quad (3.7)$$

*Then  $\bar{x}(i) \leq \bar{y}(i)$  ( $i \in \Lambda$ ).*

**Proof** Define  $\phi^n$  and  $\psi^n$  as in (3.5) and let  $x^n$  and  $y^n$  denote their unique trajectories. Then by induction, (3.7) implies that  $x^n \leq y^n$  (pointwise), so taking the limit we obtain that  $\bar{x} \leq \bar{y}$ .  $\blacksquare$

### 3.3 Complete instability

Let  $\Phi^p$  be defined as in (1.2) with  $m = 1$ , i.e.,  $\Phi^p$  is a random perturbation of the deterministic cellular automaton  $\Phi^0$  that applies the same nonconstant local monotone map  $\varphi_1 = \varphi$  in each space-time point. Let  $\bar{\rho}(p)$ , defined in (1.5), denote the density of its upper invariant law. Our Theorem 9 implies as a special case the difficult part of Toom's stability theorem (Theorem 2), which says that  $\Phi^0$  is stable if  $\varphi$  is an eroder. In the present subsection, we complement this by proving the "easy" part of Toom's stability theorem, which says that  $\Phi^0$  is completely unstable if  $\varphi$  is not an eroder.

**Lemma 29 (Complete instability)** *If  $\varphi$  is not an eroder, then  $\bar{\rho}(p) = 0$  for all  $p > 0$ .*

**Proof** By translation invariance, it suffices to prove that for each  $p > 0$ , the Markov chain  $(X_t^p)_{t \geq 0}$  defined in (1.3) and started in the initial state  $X_0^p = \underline{1}$  satisfies

$$\mathbb{P}^\perp[X_t^p(0) = 1] \xrightarrow[t \rightarrow \infty]{} 0. \quad (3.8)$$

Since  $\varphi$  is not an eroder, there exists configuration  $x \in \{0, 1\}^{\mathbb{Z}^d}$  containing finitely many zeros such that  $\Psi_\varphi^t(x) \neq \underline{1}$  for all  $t \geq 0$ . This allows us to choose for each  $t \geq 0$  a point  $i_t \in \mathbb{Z}^d$  such that  $\Psi_\varphi^t(x)(i_t) = 0$ . Let us write  $x(i) = 1 - 1_A(i)$  ( $i \in \mathbb{Z}^d$ ) where  $A \subset \mathbb{Z}^d$  is a finite set. Let  $A - i_t := \{i - i_t : i \in A\}$  ( $t \geq 0$ ). Then monotonicity implies that

$$X_t^p(0) = 0 \quad \text{a.s. on the event that} \quad \exists 0 < s \leq t \text{ s.t. } \Phi_{s,i}^p = \varphi_0 \quad \forall i \in A - i_s. \quad (3.9)$$

It follows that

$$\mathbb{P}^\perp[X_t^p(0) = 1] \leq (1 - p^{|A|})^t \quad (t \geq 0), \quad (3.10)$$

which proves (3.8).  $\blacksquare$

## 4 Construction of Toom contours

### 4.1 Minimal explanations

This section is devoted to the proofs of Theorems 23 and 26, which can be found in Subsections 4.2 and 4.3 below. In the present subsection, we prepare for these proofs by giving a formal definition of the minimal explanations that have already been mentioned several times, and investigating their properties.

Recall that  $\mathcal{A}(\phi)$  and  $\mathcal{Z}(\phi)$  denote the sets of minimal one-sets and zero-sets of a monotone local map  $\phi$ , defined in Subsections 1.1 and 1.4. In analogy with (1.9), each monotone local map  $\phi : \{0, 1\}^\Lambda \rightarrow \{0, 1\}$  can be written as

$$\phi(x) = \bigvee_{A \in \mathcal{A}(\phi)} \bigwedge_{i \in A} x(i) = \bigwedge_{Z \in \mathcal{Z}(\phi)} \bigvee_{i \in Z} x(i) \quad (x \in \{0, 1\}^\Lambda). \quad (4.1)$$



In particular, if  $\phi^0$  and  $\phi^1$  are the constant maps defined in (2.1), then  $\mathcal{Z}(\phi^0) = \{\emptyset\}$  and  $\mathcal{Z}(\phi^1) = \emptyset$ . For monotone local maps  $\phi, \phi'$ , we write

$$\phi \leq \phi' \Leftrightarrow \phi(x) \leq \phi'(x) \quad \forall x \in \{0, 1\}^\Lambda \quad \text{and} \quad \phi \preceq \phi' \Leftrightarrow \mathcal{Z}(\phi) \supset \mathcal{Z}(\phi'). \quad (4.2)$$

It is easy to see that  $\phi \preceq \phi'$  implies  $\phi \leq \phi'$ , but not the other way around. For monotone cellular automata  $\phi = (\phi_i)_{i \in \Lambda}$  and  $\phi' = (\phi'_i)_{i \in \Lambda}$ , we write  $\phi \leq \phi'$  if and only if  $\phi_i \leq \phi'_i$  for all  $i \in \Lambda$ , and similarly, we write  $\phi \preceq \phi'$  if and only if  $\phi_i \preceq \phi'_i$  for all  $i \in \Lambda$ .

**Definition 30** *Let  $\phi = (\phi_i)_{i \in \Lambda}$  be a monotone cellular automaton and let  $0 \in \Lambda$ . By definition, a minimal explanation for 0 is a monotone cellular automaton  $\phi'$  such that:*

- (i)  $\phi \preceq \phi'$  and the maximal trajectory  $\bar{x}'$  of  $\phi'$  satisfies  $\bar{x}'(0) = 0$ .
- (ii) If a monotone cellular automaton  $\phi''$  satisfies  $\phi' \preceq \phi''$  and the maximal trajectory  $\bar{x}''$  of  $\phi''$  satisfies  $\bar{x}''(0) = 0$ , then  $\phi' = \phi''$ .

**Lemma 31 (Minimal explanations)** *Let  $\phi = (\phi_i)_{i \in \Lambda}$  be a monotone cellular automaton and let  $0 \in \Lambda$ . Then the the maximal trajectory  $\bar{x}$  of  $\phi$  satisfies  $\bar{x}(0) = 0$  if and only if there exists a minimal explanation  $\phi'$  for 0.*

**Proof** If there exists a minimal explanation  $\phi'$  for 0 and  $\bar{x}$  and  $\bar{x}'$  denote the maximal trajectories of  $\phi$  and  $\phi'$ , respectively, then  $\phi \preceq \phi'$  implies  $\phi \leq \phi'$  which implies  $\bar{x} \leq \bar{x}'$  and hence in particular  $\bar{x}(0) \leq \bar{x}'(0) = 0$ . This shows that  $\bar{x}(0) = 0$  if there exists a minimal explanation for 0.

Assume, conversely, that  $\bar{x}(0) = 0$ . Let  $\Lambda_n \subset \Lambda$  be finite sets increasing to  $\Lambda$ , let  $\phi^n$  denote the monotone cellular automata defined in (3.5), and let  $x^n$  denote the unique trajectory of  $\phi^n$ . We have seen in the proof of Lemma 14 that  $\lim_{n \rightarrow \infty} x^n(0) = \bar{x}(0)$  so we can choose  $n$  large enough such that  $x^n(0) = 0$ . It is clear from the definition that  $\phi \preceq \phi^n$ . We can now step by step replace  $\phi^n$  by larger monotone cellular automata with respect to the order  $\preceq$  as long as it is possible to do so without losing the property that the trajectory is zero in 0. Since  $\mathcal{Z}(\phi_j^n) = \emptyset$  for all but finitely many  $j$  and since  $\mathcal{Z}(\phi_j^n)$  is finite for each  $j$ , this process ends after a finite number of steps, leading to a minimal explanation for 0.  $\blacksquare$

Our next proposition describes the structure of minimal explanations. In point (iii) below, we use the convention that the maximum over an empty set is zero. We call the finite directed graph  $(U, \vec{G})$  from Proposition 32 the *explanation graph* associated with the minimal explanation  $\phi'$ . The picture on the right in Figure 4 shows an example of such an explanation graph, or rather the undirected graph  $(U, G)$  associated with  $(U, \vec{G})$ .

**Proposition 32 (Explanation graphs)** *Let  $\phi = (\phi_i)_{i \in \Lambda}$  be a monotone cellular automaton and let  $(\Lambda, \vec{H})$  be its dependence graph, as defined in (2.2). Let  $0 \in \Lambda$  and let  $\phi'$  be a minimal explanation for 0. Then there exists a finite subgraph  $(U, \vec{G})$  of  $(\Lambda, \vec{H})$  with the following properties.*

- (i) *The maximal trajectory  $\bar{x}'$  of  $\phi'$  satisfies  $\bar{x}'(i) = 0$  if and only if  $i \in U$ .*
- (ii)  $\phi'_i = \phi^1$  if  $i \notin U$ .
- (iii)  $\phi'_i(x) = \bigvee_{j: (i,j) \in \vec{G}} x(j) \quad (x \in \{0, 1\}^\Lambda)$  if  $i \in U$ .
- (iv) *For each  $j \in U \setminus \{0\}$ , there exists an  $i \in U$  such that  $(i, j) \in \vec{G}$ .*

*For  $i \in U$ , the following statements are equivalent: 1.  $\phi_i = \phi^0$ , 2.  $\phi'_i = \phi^0$ , 3.  $\vec{G}_{\text{out}}(i) = \emptyset$ .*

**Proof** Define  $U := \{j \in \Lambda : \phi'_j \neq \phi^1\}$ . We claim that  $U$  is finite. This follows from the argument we have already seen in the proof of Lemma 31: if  $\Lambda_n \subset \Lambda$  are finite sets increasing to  $\Lambda$ , then for large enough  $n$  we can replace  $\phi'_j$  by  $\phi^1$  for all  $j \notin \Lambda_n$  without affecting the fact that the maximal trajectory is zero in 0. By the maximality property of  $\phi'$ , this then implies that  $\phi'_j = \phi^1$  for all  $j \notin \Lambda_n$ .

It is clear from our definition of  $U$  that  $\bar{x}'(j) = 1$  for all  $j \in \Lambda \setminus U$ . On the other hand, we cannot have  $\bar{x}'(j) = 1$  for some  $j \in U$ , since in that case we could replace  $\phi'_j$  by  $\phi^1$  while preserving the fact that the maximal trajectory is zero in 0, which contradicts the maximality property of  $\phi'$ . This proves property (i). Property (ii) is immediate from our definition of  $U$ .

We claim that for each  $i \in U$ , there exists a finite set  $Z_i \subset U$  such that  $\mathcal{Z}(\phi'_i) = \{Z_i\}$ . Indeed, property (i) and (4.1) imply that for each  $i \in U$  there exists a  $Z \in \mathcal{Z}(\phi'_i)$  such that  $\bar{x}'(k) = 0$  for all  $k \in Z$ , which by (i) implies  $Z \subset U$ . If  $\mathcal{Z}(\phi'_i)$  contains other elements apart from  $Z$ , then we can throw these away while preserving the fact that the maximal trajectory is zero in 0, contradicting the maximality property of  $\phi'$ . Now setting

$$\vec{G} := \{(i, j) : i \in U, j \in Z_i\} \quad (4.3)$$

defines a set of directed edges such that  $\vec{G} \subset \vec{H}$  and property (iii) holds. Note that in line with earlier conventions, we allow for the case that  $Z_i = \emptyset$  and  $\phi_i = \phi^0$ .

Property (iv) follows from the fact that if  $j \in U \setminus \{0\}$  and there exists no  $i \in U$  such that  $j \in Z_i$ , then by property (ii) and (4.1) we can replace  $\phi'_j$  by  $\phi^1$  while preserving the fact that the maximal trajectory is zero in all points of  $U \setminus \{j\}$ , contradicting the maximality property of  $\phi'$ .

To prove the final statement of the proposition, we observe that if  $\phi_i = \phi^0$ , then  $\vec{G}_{\text{out}}(i) \subset \vec{H}_{\text{out}}(i) = \emptyset$ , so 1. implies 3. By property (iii), 3. implies 2., which by the fact that  $\phi \prec \phi'$  in turn implies 1.  $\blacksquare$

In the special case that the monotone cellular automaton  $\phi$  is defined in terms of a typed dependence graph  $(\Lambda, \mathcal{H})$  as in Definition 21, we can strengthen Proposition 32 as follows. We call the typed directed graph  $(U, \mathcal{G})$  from the following proposition a *typed explanation graph* associated with the minimal explanation  $\phi'$ . In general,  $(U, \mathcal{G})$  is not uniquely determined by  $\phi'$ .

**Proposition 33 (Typed explanation graphs)** *Let  $(\Lambda, \mathcal{H})$  be a typed dependence graph with  $\sigma \geq 1$  types of edges and let  $\phi$  be its associated monotone cellular automaton. Let  $0 \in \Lambda$  and let  $\phi'$  be a minimal explanation for 0. Then there exists a finite typed subgraph  $(U, \mathcal{G})$  of  $(\Lambda, \mathcal{H})$  such that:*

- (i)  $\phi'_i = \phi^1$  if  $i \notin U$ ,
- (ii)  $\phi'_i = \phi^0$  if  $i \in U_* := \{i \in U : \phi_i = \phi^0\}$ ,
- (iii) for each  $i \in U \setminus U_*$  and  $1 \leq s \leq \sigma$ , there exists a  $j_s(i) \in U$  such that  $\vec{G}_{s, \text{out}}(i) = \{j_s(i)\}$ ,
- (iv)  $\phi'_i(x) = \bigvee_{s=1}^{\sigma} x(j_s(i))$  ( $x \in \{0, 1\}^\Lambda$ ) if  $i \in U \setminus U_*$ .

The untyped directed graph  $(U, \vec{G})$  associated with  $(U, \mathcal{G})$  is the explanation graph associated with the minimal explanation  $\phi'$ .

**Proof** Let  $(U, \vec{G})$  be the explanation graph associated with the minimal explanation  $\phi'$  and for each  $i \in U$ , let  $Z_i := \{j \in U : (i, j) \in \vec{G}\}$ . Then property (iii) of Proposition 32 says that

$$\phi'_i(x) = \bigvee_{j \in Z_i} x(j) \quad (i \in U, x \in \{0, 1\}^\Lambda), \quad (4.4)$$

so  $\mathcal{Z}(\phi'_i) = \{Z_i\}$  and hence  $Z_i \in \mathcal{Z}(\phi_i)$  by the fact that  $\phi_i \preceq \phi'_i$ . Let

$$A_i^s := \{j \in \Lambda : (i, j) \in \vec{H}_s\} \quad (1 \leq s \leq \sigma, i \in \Lambda). \quad (4.5)$$

Recall from Definition 21 that  $\Lambda_\bullet = \{i \in \Lambda : \phi_i \notin \{\phi^0, \phi^1\}\}$  and that

$$\phi_i(x) = \bigvee_{s=1}^{\sigma} \bigwedge_{j \in A_i^s} x(j) \quad (i \in U, x \in \{0, 1\}^\Lambda). \quad (4.6)$$

We claim that

$$A_i^s \cap Z \neq \emptyset \quad (Z \in \mathcal{Z}(\phi_i), 1 \leq s \leq \sigma, i \in \Lambda_\bullet). \quad (4.7)$$

Indeed, if we would have  $A_i^s \cap Z = \emptyset$  for some  $Z \in \mathcal{Z}(\phi_i)$ ,  $1 \leq s \leq \sigma$ , and  $i \in \Lambda_\bullet$ , then  $1 = \varphi_i(1_{A_i^s}) \leq \varphi_i(1 - 1_Z) = 0$ , which is a contradiction. By (4.7), for each  $1 \leq s \leq \sigma$  and  $i \in \Lambda_\bullet$ , we can choose  $j_s(i) \in A_i^s \cap Z_i$ . Let  $Z'_i := \{j_s(i) : 1 \leq s \leq \sigma\}$ . Then clearly  $Z'_i \subset Z_i$ . We claim that in fact  $Z'_i = Z_i$ . Indeed, since  $j_s(i) \in A_i^s$  ( $1 \leq s \leq \sigma$ ), formula (4.6) shows that  $\phi_i(1 - 1_{Z'_i}) = 0$ . Since  $Z'_i \subset Z_i$ , by the minimality of the latter, we conclude that  $Z'_i = Z_i$ . As a result, defining a typed directed graph  $(U, \mathcal{G})$  with  $\sigma$  types of edges by setting

$$\vec{G}_s := \{(i, j_s(i)) : i \in U \cap \Lambda_\bullet\} \quad (1 \leq s \leq \sigma), \quad (4.8)$$

we have that  $(U, \vec{G})$  is the untyped directed graph associated with  $(U, \mathcal{G})$  and

$$\phi'_i(x) = \bigvee_{j \in Z'_i} x(j) = \bigvee_{s=1}^{\sigma} x(j_s(i)) \quad (i \in U \cap \Lambda_\bullet, x \in \{0, 1\}^\Lambda). \quad (4.9)$$

Now properties (i)–(iv) follow from properties (ii) and (iii) of Proposition 33, while (4.5) shows that  $(U, \mathcal{G})$  is a typed subgraph of  $(\Lambda, \mathcal{H})$ .  $\blacksquare$

## 4.2 Toom contours

In this subsection, we prove Theorem 23. We fix a typed dependence graph  $(\Lambda, \mathcal{H})$  with  $\sigma \geq 1$  types of edges. We let  $\phi$  denote its associated monotone cellular automaton and let  $\bar{x}$  denote its maximal trajectory. We fix an element  $0 \in \Lambda$  and assume that  $\bar{x}(0) = 0$ . We need to prove the presence in  $(\Lambda, \mathcal{H})$  of a Toom contour  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  rooted at  $0$ . Since  $\bar{x}(0) = 0$ , by Lemma 31, there exists a minimal explanation  $\phi'$  for  $0$ , and by Proposition 33, there exists a typed explanation graph  $(U, \mathcal{G})$  associated with  $\phi'$ . We will derive Theorem 23 from the following theorem. Recall Definition 18 of an embedding of a rooted Toom graph.

**Theorem 34 (Toom graph embedded in explanation graph)** *Let  $(U, \mathcal{G})$  be a typed explanation graph associated with a minimal explanation  $\phi'$  for  $0$ . Then there exists a rooted Toom graph  $(v_\circ, \mathcal{V}, \mathcal{E})$  and an embedding  $\psi$  of  $(v_\circ, \mathcal{V}, \mathcal{E})$  in  $U$  such that  $\psi(v_\circ) = 0$  and*

- (i)  $\psi(V_\ast) = U_\ast$ ,
- (ii)  $(\psi(v), \psi(w)) \in \vec{G}_s$  for all  $(v, w) \in \vec{E}_s^\bullet$  ( $1 \leq s \leq \sigma$ ),
- (iii)  $(\psi(v), \psi(w)) \in \vec{G}$  for all  $(v, w) \in \vec{E}^\circ$ ,

where  $\vec{E}_s^\bullet$  and  $\vec{E}^\circ$  are defined in (2.15).

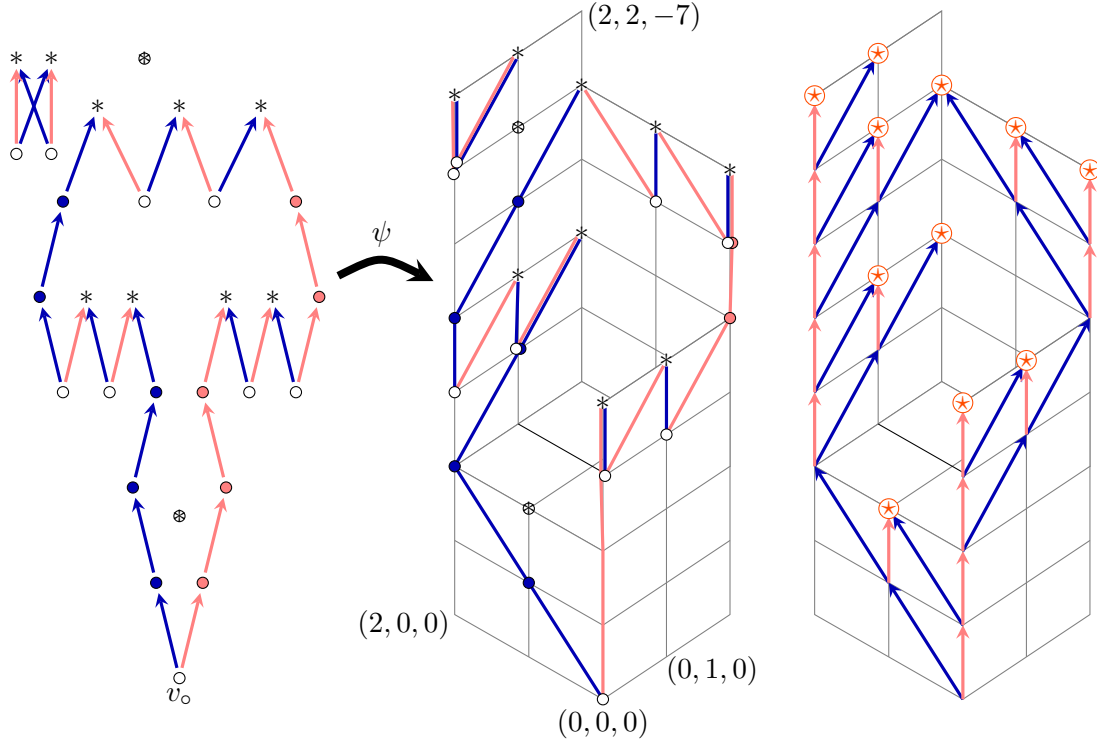


Figure 4: Embedding of a rooted Toom graph inside a typed explanation graph. On the right: a typed explanation graph  $(U, \mathcal{G})$  associated with a minimal explanation for  $(0, 0, 0)$  in the sense of Proposition 33. On the left and in the middle: embedding of a rooted Toom graph in  $(U, \mathcal{G})$  in the sense of Theorem 34. The connected component of this Toom graph containing the root is a Toom contour rooted at  $(0, 0, 0)$  (compare Figure 3).

To see that Theorem 34 implies Theorem 23, it suffices to observe that if  $(V', \mathcal{E}')$  is the connected component containing  $v_o$  of the Toom graph  $(\mathcal{V}, \mathcal{E})$  from Theorem 34, and  $\psi'$  is the restriction of  $\psi$  to  $V'$ , then  $(v_o, V', \mathcal{E}', \psi')$  is a Toom contour rooted at 0 that is present in  $(\Lambda, \mathcal{H})$ . Note that when we restrict ourselves to the connected component containing the root, property (i) of Theorem 34 must be weakened to  $\psi(V_*) \subset U_*$ , which is all that is needed to satisfy Definition 22 (i). Theorem 34 is demonstrated in Figure 4.

We observe that Theorem 34 is trivial if  $|U| = 1$ , since in this case 0 is a defective site and we can take for  $(\mathcal{V}, \mathcal{E})$  the trivial Toom graph that consists of a single isolated vertex. We assume therefore from now on that  $|U| \geq 2$ . In this case,  $0 \notin U_*$ .

The proof of Theorem 34 needs some preparations. In any directed graph  $(V, \vec{E})$ , for two vertices  $i, j \in V$ , we write  $i \rightsquigarrow j$  if there exists  $i = i_0, \dots, i_n = j$  such that  $(i_{k-1}, i_k) \in \vec{E}$  ( $1 \leq k \leq n$ ). By definition, a *time-ordering* of  $(V, \vec{E})$  is an enumeration  $V = \{i_1, \dots, i_N\}$  of its vertices such that for each  $1 \leq n \leq N$ , there are no  $k, l$  with  $k < n \leq l$  and  $(i_l, i_k) \in \vec{E}$ . Note that since  $(\Lambda, \vec{H})$  is acyclic, the same is true for  $(U, \vec{G})$ .

**Lemma 35 (Time-ordering)** *Each finite acyclic directed graph has a time-ordering. For an explanation graph  $(U, \vec{G})$ , we can choose a time-ordering such that  $i_1 = 0$  and  $U \setminus U_* = \{i_1, \dots, i_m\}$  for some  $1 \leq m \leq |U|$ .*

**Proof** For any acyclic directed graph, the relation  $\rightsquigarrow$  is a partial order on  $V$ ; in particular, there cannot exist  $i, j \in V$  with  $i \neq j$  such that  $i \rightsquigarrow j$  and  $j \rightsquigarrow i$ , since this would imply the existence of a cycle in  $(V, \vec{E})$ . We can now inductively construct a time-ordering  $i_1, i_2, \dots$  by choosing for  $i_n$  a minimal element of  $V \setminus \{i_1, \dots, i_{n-1}\}$ .

It follows from Proposition 32 (iv) that  $0 \rightsquigarrow i$  for all  $i \in U$ , so by the fact that  $(U, \vec{G})$  is acyclic we have  $i \not\rightsquigarrow 0$  for all  $i \in U \setminus \{0\}$ . Thus, 0 is a minimal element of  $U$  with respect to the partial order  $\rightsquigarrow$  and we can construct the time-ordering starting with  $i_1 = 0$ . Since elements of  $U_*$  have no outgoing edges, we can also first construct a time-ordering of  $U \setminus U_*$  and then add the elements of  $U_*$  in any order.  $\blacksquare$

From now on, we fix a typed explanation graph  $(U, \mathcal{G})$  associated with a minimal explanation  $\phi'$  for 0, as well as a time-ordering of the associated untyped explanation graph  $(U, \vec{G})$  with the properties described in Lemma 35. We let  $m := |U \setminus U_*|$  and we adopt the following definitions.

**Definition 36** For each  $1 \leq n \leq m$ , we set  $U_n^- := \{i_1, \dots, i_n\}$  and  $U_n^+ := U \setminus U_n^-$ . We call

$$\partial U_n^- := \{j \in U_n^+ : \exists i \in U_n^- \text{ s.t. } (i, j) \in \vec{G}\} \quad (4.10)$$

the boundary of  $U_n^-$ . We equip  $\partial U_n^-$  with the structure of an unoriented graph in which two elements  $i, j \in \partial U_n^-$  are neighbours, denoted  $i \approx j$ , if there exists a  $k \in U$  such that  $i \rightsquigarrow k$  and  $j \rightsquigarrow k$ . We write  $i \sim j$  if  $i, j \in \partial U_n^-$  lie in the same connected component of this graph.

The following lemma says that the number of connected components on the boundary  $\partial U_n^-$  is non-decreasing in  $n$ . Note that at the end, when  $n = m$ , we have  $\partial U_m^- = U_*$  and each element of  $\partial U_m^-$  forms a connected component on its own. Therefore, starting from a single connected component at  $n = 1$  the boundary gradually breaks up into smaller and smaller connected components.

**Lemma 37 (Break-up of boundary)** For each  $1 < n \leq m$ , if  $C$  is a connected component of  $\partial U_{n-1}^-$  and  $i_n \notin C$ , then  $C$  is also a connected component of  $\partial U_n^-$ . Each connected component of  $\partial U_n^-$  that is not a connected component of  $\partial U_{n-1}^-$  contains a vertex  $j$  such that  $(i_n, j) \in \vec{G}$ .

**Proof** We first prove that a connected component  $C$  of  $\partial U_{n-1}^-$  that does not contain  $i_n$  is also a connected component of  $\partial U_n^-$ . For each  $i, j \in C$ , there exist  $i(0), \dots, i(k) \in C$  such that  $i(0) \approx \dots \approx i(k)$  with  $i(0) = i$  and  $i(k) = j$ , which implies that  $i \sim j$  in  $U_n^-$ . This shows that  $C$  is contained in some connected component  $C'$  of  $\partial U_n^-$ . We need to show that  $C = C'$ . Assume that conversely,  $C'$  is strictly larger than  $C$ . Then we can find  $i \in C$  and  $j \in C' \setminus C$  such that  $i \approx j$ . Since  $j \in \partial U_n^-$  we must have either  $j \in \partial U_{n-1}^-$  or  $(i_n, j) \in \vec{G}$  (possibly both). If  $j \in \partial U_{n-1}^-$ , then  $i \approx j$  implies  $j \in C$  which contradicts our assumptions. However, if  $(i_n, j) \in \vec{G}$ , then  $i \approx j$  implies  $i \approx i_n$  which also contradicts our assumptions, since  $C$  does not contain  $i_n$ .

To prove the second claim of the lemma, assume that  $C$  is a connected component of  $\partial U_n^-$  that is not a connected component of  $\partial U_{n-1}^-$ . Let  $i$  be any element of  $C$ . If  $(i_n, i) \in \vec{G}$  we are done. In the opposite case,  $i \in \partial U_{n-1}^-$ . Since  $C$  is not a connected component of  $\partial U_{n-1}^-$ , by what we have already proved,  $i$  must lie in the connected component of  $\partial U_{n-1}^-$  that contains  $i_n$ , so there exist  $i(0), \dots, i(k) \in \partial U_{n-1}^-$  with  $i(0) = i$ ,  $i(k) = i_n$ , and  $i(0) \approx \dots \approx i(k)$ . Now there exists a  $j' \in U$  such that  $i(k-1) \rightsquigarrow j'$  and  $i_n \rightsquigarrow j'$ . If  $j' \neq i_n$ , then let  $j$  be the first vertex after  $i_n$  on the path from  $i_n$  to  $j'$ , and if  $j' = i_n$ , then choose for  $j$  any vertex with  $(i_n, j) \in \vec{G}$ . In either case, we then have  $i(k-1) \approx j$  which implies that  $j \in C$ , and clearly  $(i_n, j) \in \vec{G}$ .  $\blacksquare$

**Definition 38** For  $1 \leq s \leq \sigma$ , we define a spoke of charge  $s$  to be a sequence  $(i(0), \dots, i(k))$  of vertices in  $U$  such that  $k \geq 1$ ,  $i(k) \in U_*$ ,  $(i(0), i(1)) \in \vec{G}$ , and  $(i(l-1), i(l)) \in \vec{G}_s$  for all  $2 \leq l \leq k$ . We say that a spoke  $(i(0), \dots, i(k))$  intersects a set  $V \subset U$  if  $i(l) \in V$  for some  $0 \leq l \leq k$ . A pole at vertex  $i \in U$  is a collection  $(i_s(0), \dots, i_s(k_s))_{1 \leq s \leq \sigma}$  of spokes of charges  $1 \leq s \leq \sigma$ , respectively, such that  $i_s(0) = i$  for all  $1 \leq s \leq \sigma$ .

**Proof of Theorem 34** We have already shown that the statement is trivial if 0 is a defective site, so we continue assuming that  $|U| > 1$  and  $0 \notin U_*$ . We fix a time-ordering of  $(U, \vec{G})$  as in Lemma 35 and define  $U_n^\pm$  as in Definition 36. Let  $N(n)$  denote the number of connected components of  $\partial U_n^-$  ( $1 \leq n \leq m$ ). It follows from Lemma 37 that  $N(n)$  increases to  $N(m) = |U_*|$ . We will show by induction that for each  $1 \leq n \leq m$ , it is possible to construct poles

$$(i_s^r(0), \dots, i_s^r(k_s^r))_{1 \leq s \leq \sigma} \quad (1 \leq r \leq N(n)) \quad (4.11)$$

at vertices  $i^1, \dots, i^{N(n)} \in U_n^-$  such that

- (i)  $i^1 = 0$  and  $(i_s^1(0), i_s^1(1)) \in \vec{G}_s$  ( $1 \leq s \leq \sigma$ ),
- (ii) for each connected component  $C$  of  $\partial U_n^-$  and for each  $1 \leq s \leq \sigma$ , there exists precisely one  $1 \leq r \leq N(n)$  such that the spoke  $(i_s^r(0), \dots, i_s^r(k_s^r))$  intersects  $C$ .

We start by proving the claim for  $n = 1$ . By Proposition 33 (iii), for each  $1 \leq s \leq \sigma$ , at each  $i \in U \setminus U_*$  there is precisely one outgoing edge of charge  $s$ . Thus, for each  $1 \leq s \leq \sigma$ , there starts a unique spoke  $(i_s^1(0), \dots, i_s^1(k_s^1))$  at 0 such that  $(i_s^1(l-1), i_s^1(l)) \in \vec{G}_s$  for all  $1 \leq l \leq k_s$ , and these spokes together form a pole at 0 such that (i) holds. If  $\partial U_1^-$  has only one connected component, then (ii) also holds and we are now done. In the opposite case, we can add additional poles at 0 so that (ii) holds.

We now continue by induction on  $n$ . We will show that by adding poles, we can make sure (ii) remains valid as we increase  $n$ . Since (i) also obviously stays true if we add poles, this then completes the proof that (i) and (ii) can be satisfied for all  $n$ . Assume that we have poles at vertices  $i^1, \dots, i^{N(n-1)} \in U_{n-1}^-$  such that conditions (i) and (ii) are satisfied. By Lemma 37, if  $C$  is a connected component of  $\partial U_{n-1}^-$  that does not contain  $i_n$ , then  $C$  is also a connected component of  $\partial U_n^-$ , so for such a connected component  $C$  condition (ii) remains satisfied even without adding new poles. Let  $C_1, \dots, C_k$  be the other connected components of  $\partial U_{n-1}^-$ , which contain all vertices of the connected component of  $\partial U_{n-1}^-$  that contains  $i_n$ , except  $i_n$  itself, as well as all vertices  $j \in U$  such that  $(i_n, j) \in \vec{G}$ . Using this and the induction hypothesis, we see that for each  $1 \leq s \leq \sigma$ , there exists precisely one  $1 \leq r \leq N(n-1)$  such that the spoke  $(i_s^r(0), \dots, i_s^r(k_s^r))$  intersects  $C_1 \cup \dots \cup C_k$ . By Lemma 37, each of the components  $C_1, \dots, C_k$  contains an element  $j$  such that  $(i_n, j) \in \vec{G}$ . Using this and the fact that the edge between the first two vertices along each pole is in  $\vec{G}$ , we see that we can add additional poles in  $i_n$  so that condition (ii) is satisfied for  $C_1, \dots, C_k$ . This completes the induction step.

In particular, setting  $n = m$ , we have now shown that it is possible to construct poles

$$(i_s^r(0), \dots, i_s^r(k_s^r))_{1 \leq s \leq \sigma} \quad (1 \leq r \leq |U_*|) \quad (4.12)$$

at vertices in  $U \setminus U_*$  such that

- (i)  $i^1 = 0$  and  $(i_s^1(0), i_s^1(1)) \in \vec{G}_s$  ( $1 \leq s \leq \sigma$ ),
- (ii) for each  $i \in U_*$  and for each  $1 \leq s \leq \sigma$ , there exists precisely one  $1 \leq r \leq |U_*|$  such that the spoke  $(i_s^r(0), \dots, i_s^r(k_s^r))$  ends in  $i_s^r(k_s^r) = i$ .

It is straightforward to check that these poles together define a Toom graph  $(\mathcal{V}, \mathcal{E})$  that is embedded in  $(U, \mathcal{G})$  in such a way that conditions (i)–(iii) of the theorem are satisfied. Indeed, each pole corresponds to a source and its  $\sigma$  spokes to the  $\sigma$  charges emerging from the source. To see that  $\psi$  satisfies condition (ii) of Definition 18 of an embedding of a rooted Toom graph, one uses the fact that if two spokes of the same charge would enter the same vertex, then these spokes would have to be equal starting from that vertex, which would lead to two spokes of the same charge ending in the same defective site, contradicting point (ii) above. Also, since there are no incoming edges at 0 in the explanation graph, we never add internal vertices that

overlap with the root. The fact that conditions (ii) and (iii) of the theorem are satisfied follows our definition of a spoke of charge  $s$ , which includes the condition that  $(i(l-1), i(l)) \in \vec{G}_s$  for all  $2 \leq l \leq k$ , as well as point (i) above.  $\blacksquare$

### 4.3 Toom cycles

In this subsection, we prove Theorem 26. As in the proof of Theorem 23, we will construct the Toom cycle inside a typed explanation graph. Apart from this similarity, the proof will be completely different.

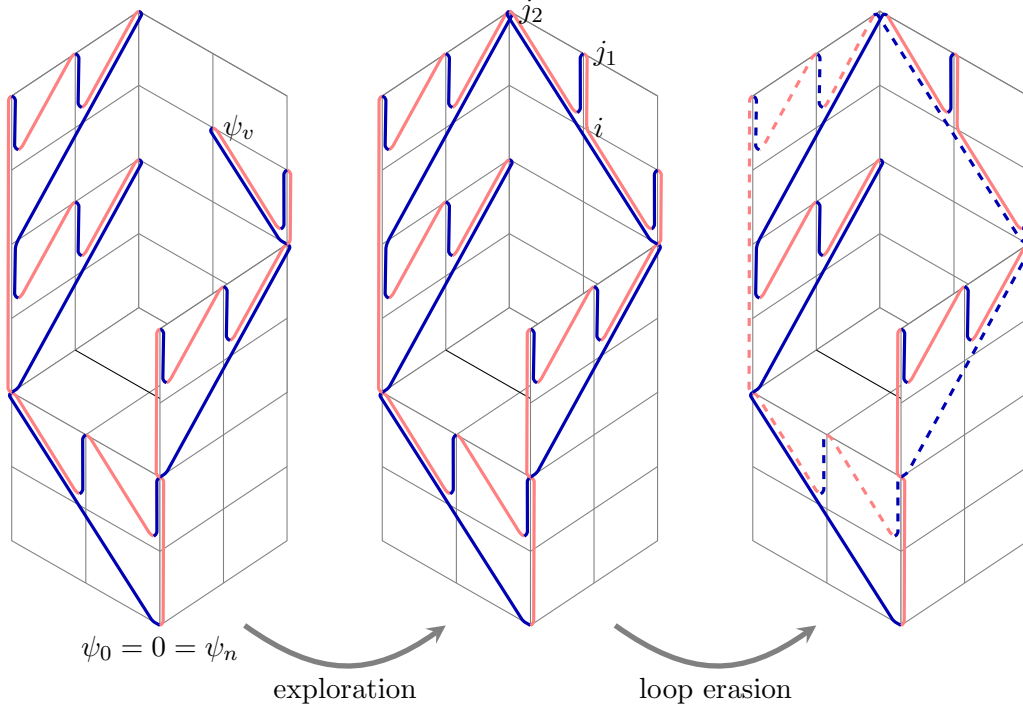


Figure 5: The process of exploration and loop erasure. The Toom cycle is constructed on the explanation graph of Figure 4. We can see that in the Toom cycle on the left  $v$  is a sink, but  $\psi_v = i$  is not a defective site. In the exploration step,  $v$  is replaced by two internal vertices, one of each charge, and two new sinks are added to the cycle at the positions  $j_1$  and  $j_2$ . This leads to the new sink at  $j_2$  overlapping with a preexisting sink. In the loop erasure step, this is resolved by erasing the part of the cycle between the first and second visit to  $j_2$ .

**Proof of Theorem 26** We fix a typed explanation graph  $(U, \mathcal{G})$  as in Proposition 33, with two types of edges. By Proposition 32 (iv), for each  $i \in U$ , there exist  $i_0, \dots, i_n \in U$  with  $i_0 = 0$  and  $i_n = i$  such that  $(i_{k-1}, i_k) \in \vec{G}$  for all  $1 \leq k \leq n$ . We let  $\text{dist}(i)$  denote the smallest integer  $n$  for which such  $i_0, \dots, i_n$  can be found, i.e.,  $\text{dist}(i)$  is the length of the shortest directed path in  $(U, \vec{G})$  from 0 to  $i$ . We will use an inductive construction. At each point in the construction, we have a Toom cycle  $(V, \vec{E}, \psi)$  rooted at 0, and we let

$$M := \sup_{v \in V} \text{dist}(\psi_v) \tag{4.13}$$

denote the largest distance from 0 of all vertices of the Toom cycle. We will make sure that at each point in our construction, the following induction hypotheses are satisfied:

- (i)' if  $\psi_v \notin U_*$  for some  $v \in V_*$ , then  $\text{dist}(\psi_v) \geq M - 1$ ,

- (ii)  $(\psi_v, \psi_w) \in \vec{H}_s$  for all  $(v, w) \in \vec{E}_s$  with  $v \in V'_s$  ( $s = 1, 2$ ),
- (iii)  $(\psi_v, \psi_w) \in \vec{H}_{3-s}$  for all  $(v, w) \in \vec{E}_s$  with  $v \in V'_o$  ( $s = 1, 2$ ).

Our construction will end as soon as in place of (i)' we have the stronger condition

- (i)  $\psi_v \in U_*$  for all  $v \in V_*$ ,

since this then guarantees that  $(V, \vec{E}, \psi)$  is present in  $(\Lambda, \mathcal{H})$ . We note that in order to specify the Toom cycle, it suffices to know the function  $\psi : [n] \rightarrow U$  only, since the induction hypotheses (ii) and (iii) imply that for  $1 \leq v \leq n$ ,

$$(v-1, v) \in \vec{E} \text{ if } (\psi_{v-1}, \psi_v) \in \vec{H} \quad \text{and} \quad (v, v-1) \in \vec{E} \text{ if } (\psi_v, \psi_{v-1}) \in \vec{H}. \quad (4.14)$$

Thus, at each step in the induction, we only specify an integer  $n \geq 1$  and a function  $\psi : [n] \rightarrow U$ ; it is then implicit that  $\vec{E}$  is defined by (4.14). It will be useful to view  $\psi$  as a word  $\psi_0 \cdots \psi_n$  of length  $n+1$ , made up from the alphabet  $U$ , with  $\psi_0 = 0 = \psi_n$ . We start with  $n = 1$  and  $\psi_0 = \psi_1 := 0$ . If  $0 \in U_*$  (and hence  $|U| = 1$ ), then we are done. In the opposite case, as long as (i) is not yet satisfied, we modify  $\psi$  according to the following two steps, that are illustrated in Figure 5.

- I. *Exploration.* We pick a  $v \in V_*$  such that  $i := \psi_v \notin U_*$ . If it is possible to pick  $v$  such that  $\text{dist}(\psi_v) = M-1$ , then we do so; in the opposite case we pick  $v$  such that  $\text{dist}(\psi_v) = M$ . By Proposition 33 (iii) there are unique  $j_1, j_2 \in U$  such that  $(i, j_s) \in \vec{G}_s$  ( $s = 1, 2$ ). We modify the word  $\psi_0 \cdots \psi_n$  by inserting in place of the letter  $\psi_v = i$  the string  $i j_1 i j_2 i$ .
- II. *Loop erasion.* If as the result of the exploration, there are  $v_1, v_2 \in V_*$  with  $v_1 < v_2$  such that  $i := \psi_{v_1} = \psi_{v_2}$ , then in place of the string  $\psi_{v_1} \cdots \psi_{v_2}$  we insert the letter  $i$ . We repeat this until there are no more  $v_1, v_2 \in V_*$  with  $v_1 < v_2$  such that  $\psi_{v_1} = \psi_{v_2}$ .

We must check that at the end of each induction step, we obtain a Toom cycle satisfying the induction hypotheses (i)', (ii), (iii). We first investigate the effect of exploration.

After the exploration step, it is clear that  $\psi$  via (4.14) defines an oriented cycle. The map  $\psi$  may no longer satisfy condition (i) of Definition 24 (this will be fixed in the loop erasion step), but because of the way we have chosen  $v$ , after the exploration step, it will be true that:

$$\begin{aligned} &\text{if } \psi_w = \psi_{w'} \text{ for some } w \in V_* \text{ and } w' \in V \text{ with } w \neq w', \\ &\text{then } \text{dist}(w) = \text{dist}(w') = M \text{ and } w' \in V_*. \end{aligned} \quad (4.15)$$

Indeed, the fact that  $\text{dist}(w) = \text{dist}(w') = M$  follows from the fact that the newly added vertices are at the largest distance  $M$  from 0, while (2.17) and (4.14) imply that vertices at distance  $M$  from 0 must be elements of  $V_*$ . We claim that after the exploration step  $\psi$  still satisfies condition (ii) of Definition 24. Indeed, before the exploration step, (i) was still satisfied so the sink at  $i$  did not overlap with any other vertices. In the exploration step, we add two sinks at the positions  $j_1$  and  $j_2$  and replace the old sink at the position  $i$  by three vertices in  $V_1, V_o$ , and  $V_2$ , respectively, in this order. From this we see that after the exploration step, condition (ii) of Definition 24 is still satisfied.

We claim that after the exploration step, the induction hypotheses (i)', (ii), and (iii) remain valid. Indeed, (i)' remains valid since  $M$  does not increase unless all  $v \in V_*$  for which  $\psi_v \notin U_*$  are at distance  $M$  from 0, and hence at least at distance  $M-1$  after  $M$  has increased. For the remaining induction hypotheses, we observe that in the exploration step, all existing edges of the oriented cycle keep their orientation. Their starting and endvertices also stay in whichever of the sets  $V_o, V_*, V_1$ , and  $V_2$  they were in before, except (in the case  $v \neq 0$ ) for the edges that ended in the vertex  $v \in V_*$ , whose new endvertices now belong to the sets  $V_1$



and  $V_2$ , respectively. This has no influence on the induction hypotheses (ii) and (iii), however, for which only the starting vertices matter. Also, it is straightforward to check that the new edges inserted in the exploration step satisfy (ii) and (iii).

We next investigate the effect of loop erasion. During a loop erasion, all edges keep their charge and (because we are assuming  $v_1, v_2 \in V_*$ ) also all vertices stay in whichever of the sets  $V_0, V_*, V_1$ , and  $V_2$  they were in before. Since they moreover preserve their relative order in  $V$ , this implies that condition (ii) of Definition 24 and the induction hypotheses (ii) and (iii) remain valid. In view of (4.15), during loop erasion,  $M$  does not change and hence the induction hypothesis (i)' also remains valid. Furthermore, the process of loop erasion also restores condition (i) of Definition 24. This completes the induction step.

To complete the proof, we observe that in each step, either  $M$  increases, or the number of vertices  $v \in V_*$  with  $\text{dist}(\psi_v) = M - 1$  and  $\psi_v \notin U_*$  decreases. By the finiteness of  $(U, \mathcal{G})$ , this implies that our inductive construction terminates after a finite number of steps. It follows from our induction hypotheses and the fact that  $(U, \mathcal{G})$  is a subgraph of the typed dependence graph  $(\Lambda, \mathcal{H})$  that at the end we obtain a Toom cycle that is present in  $(\Lambda, \mathcal{H})$ .  $\blacksquare$

## 5 The Peierls argument

### 5.1 Set-up

In this section we prove Theorem 9, which gives sufficient conditions for the stability of monotone cellular automata with intrinsic randomness, as well as Proposition 11, which gives lower bounds on the critical noise parameter for two deterministic monotone cellular automata. Both proofs are based on the Peierls bound from Theorem 27. In the present subsection we translate the setting of Theorem 9 and Proposition 11 into the more general language of Theorem 27 and make a choice for the typed dependence graph  $(\Lambda, \mathcal{H})$  of Theorem 27 based on a linear polar function.

Throughout this section we assume that:

- $\varphi_1, \dots, \varphi_m : \{0, 1\}^{\mathbb{Z}^d} \rightarrow \{0, 1\}$  are non-constant monotone local functions,
- $\mathbf{r} = (\mathbf{r}(1), \dots, \mathbf{r}(m))$  is a probability distribution on  $\{1, \dots, m\}$ .

For each  $p \in [0, 1]$ , we let  $\Phi^p = (\Phi_{i,t}^p)_{(i,t) \in \mathbb{Z}^{d+1}}$  be an i.i.d. collection of maps as in (1.2). We set  $\Lambda := \mathbb{Z}^{d+1}$  and for each  $(i, t) \in \Lambda$  (with  $i \in \mathbb{Z}^d$  and  $t \in \mathbb{Z}$ ) we define  $\Phi_{(i,t)}^p$  as in (2.4), so that  $(\Phi_{(i,t)}^p)_{(i,t) \in \Lambda}$  is a random monotone cellular automaton of the type considered in Theorem 27. It will be convenient to define  $\kappa : \Lambda \rightarrow \{0, \dots, m\}$  by

$$\Phi_{i,t} =: \varphi_{\kappa(i,t)} \quad ((i, t) \in \Lambda). \quad (5.1)$$

Then, in the notation of Theorem 27,

$$\Lambda_0 = \{(i, t) \in \Lambda : \kappa(i, t) = 0\} \quad \text{and} \quad \Lambda_\bullet = \{(i, t) \in \Lambda : \kappa(i, t) \in \{1, \dots, m\}\}. \quad (5.2)$$

In order to apply Theorem 27, we need to choose  $A_{s,(i,t)} \in \mathcal{A}(\Phi_{(i,t)})$  for each  $1 \leq s \leq \sigma$  and  $(i, t) \in \Lambda_\bullet$ . We will let our choice be guided by a polar function. Throughout this section we assume that  $L : \mathbb{R}^d \rightarrow \mathbb{R}^\sigma$  is a linear polar function of dimension  $\sigma \geq 2$  such that

$$\varepsilon := \sum_{s=1}^{\sigma} \varepsilon_s > 0 \quad \text{with} \quad \varepsilon_s := \inf_{1 \leq k \leq m} \varepsilon_{\varphi_k}(L_s) \quad (1 \leq s \leq \sigma), \quad (5.3)$$

where  $\varepsilon_{\varphi_k}(L_s)$  is the edge speed defined in (1.15). For each  $1 \leq s \leq \sigma$  and  $1 \leq k \leq m$ , we fix  $A_{s,k} \in \mathcal{A}(\varphi_k)$  such that

$$\sup_{A \in \mathcal{A}(\varphi_k)} \inf_{i \in A} L_s(i) =: \varepsilon_{\varphi_k}(L_s) = \inf_{i \in A_{s,k}} L_s(i) \quad (1 \leq s \leq \sigma, 1 \leq k \leq m), \quad (5.4)$$

i.e.,  $A_{s,k}$  is a set for which the supremum in the definition of the edge speed in (1.15) is achieved. Then setting

$$A_{s,(i,t)} := \{(i+j, t-1) : j \in A_{s,\kappa(i,t)}\} \quad (1 \leq s \leq \sigma, (i,t) \in \Lambda_\bullet) \quad (5.5)$$

defines sets  $A_{s,(i,t)} \in \mathcal{A}(\Phi_{(i,t)})$  as needed for the application of Theorem 27. In Subsection 5.3 below we will explain why this is the “right” choice for these sets. We let  $(\Lambda, \mathcal{H})$  denote the typed dependence graph defined in terms of the sets  $A_{s,(i,t)}$  as in (2.21), and  $(\overline{X^p}(i,t))_{(i,t) \in \Lambda}$  denote the maximal trajectory of  $\Phi^p$ . We denote the origin in  $\mathbb{Z}^{d+1} = \mathbb{Z}^d \times \mathbb{Z}$  by  $(0,0)$ , and write  $\mathcal{T}_{(0,0)}$  for the set of Toom contours rooted at  $(0,0)$ , up to equivalence. Then Theorem 27 tells us that

$$1 - \bar{\rho}(p) = \mathbb{P}[\overline{X^p}(0,0) = 0] \leq \sum_{T \in \mathcal{T}_{(0,0)}} \mathbb{P}[T \text{ is present in } (\Lambda, \mathcal{H})]. \quad (5.6)$$

We observe that as a consequence of properties (ii) and (iii) of Definition 22, each Toom contour  $T = (v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  with  $\sigma$  charges that is present in  $(\Lambda, \mathcal{H})$  must satisfy:

$$\begin{aligned} \text{(ii)'} \quad & \psi(w) = \psi(v) + (j, -1) \text{ for some } j \in \Delta_s \text{ for all } (v,w) \in \vec{E}_s^\bullet \quad (1 \leq s \leq \sigma), \\ \text{(iii)'} \quad & \psi(w) = \psi(v) + (j, -1) \text{ for some } j \in \Delta \text{ for all } (v,w) \in \vec{E}^\circ, \end{aligned} \quad (5.7)$$

where

$$\Delta_s := \bigcup_{k=1}^m A_{s,k} \quad (1 \leq s \leq \sigma) \quad \text{and} \quad \Delta := \bigcup_{s=1}^\sigma \Delta_s. \quad (5.8)$$

We let  $\mathcal{T}'_{(0,0)}$  denote the set of all  $T \in \mathcal{T}_{(0,0)}$  that satisfy (5.7). Then clearly, in (5.6) we can restrict the sum to  $T \in \mathcal{T}'_{(0,0)}$  since all other terms are zero.

For Toom cycles, similar arguments apply. In this case, we won't need the concept of equivalence of Toom contours defined in Definition 20 but can use the slightly weaker but more intuitive concept of isomorphism of Toom contours. In line with this, we let  $\vec{\mathcal{T}}_{(0,0)}$  denote the set of Toom cycles rooted at  $(0,0)$ , up to isomorphism, and inspired by Definition 25, we let  $\vec{\mathcal{T}}'_{(0,0)}$  denote the subset of Toom cycles that moreover satisfy, for  $s = 1, 2$ ,

$$\begin{aligned} \text{(ii)''} \quad & \psi(w) = \psi(v) + (j, -1) \text{ for some } j \in \Delta_s \text{ for all } (v,w) \in \vec{E}_s \text{ with } v \in V'_s, \\ \text{(iii)''} \quad & \psi(w) = \psi(v) + (j, -1) \text{ for some } j \in \Delta_{3-s} \text{ for all } (v,w) \in \vec{E}_s \text{ with } v \in V'_\circ. \end{aligned} \quad (5.9)$$

Then Theorem 27 tells us that

$$1 - \bar{\rho}(p) = \mathbb{P}[\overline{X^p}(0,0) = 0] \leq \sum_{T \in \mathcal{T}'_{(0,0)}} \mathbb{P}[T \text{ is present in } (\Lambda, \mathcal{H})]. \quad (5.10)$$

## 5.2 Stability of cellular automata with intrinsic randomness

In this subsection we prove Theorem 9. For each Toom contour  $T = (v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  rooted at  $(0,0)$  let

$$n_*(T) := |V_\circ| = |V_*| \quad (5.11)$$

denote its number of sinks and sources, each. Recall that  $\mathcal{T}'_{(0,0)}$  denotes the set of Toom contours with the additional properties (ii)' and (iii)' from (5.7). The following lemma states that each  $T \in \mathcal{T}'_{(0,0)}$  has an equal number of charged edges of each charge.

**Lemma 39 (Number of charged edges)** *For each Toom contour  $T = (v_\circ, \mathcal{V}, \mathcal{E}, \psi) \in \mathcal{T}'_{(0,0)}$  with  $\sigma \geq 2$  charges there exists an integer  $n_e(T)$  such that*

$$n_e(T) := |\vec{E}_1| = \dots = |\vec{E}_\sigma|. \quad (5.12)$$

**Proof** We write  $\psi(v) = (\psi_1(v), \dots, \psi_{d+1}(v))$  where  $\psi_{d+1}(v)$  denotes the time coordinate. The conditions in (5.7) imply that  $\psi_{d+1}(v) - \psi_{d+1}(w) = 1$  for each  $(v, w) \in \vec{E}$ . Recall that by Definition 16 in a Toom graph at each source there emerge  $\sigma$  charges, one of each type, that then travel via internal vertices of the corresponding charge through the graph until they arrive at a sink, in such a way that at each sink there converge precisely  $\sigma$  charges. This implies

$$|\vec{E}_1| = \dots = |\vec{E}_\sigma| = \sum_{v \in V_*} \psi_{d+1}(v) - \sum_{v \in V_o} \psi_{d+1}(v). \quad (5.13)$$

■

To prove Theorem 9 we need two more lemmas, the proof of which will be postponed till later. To state the first lemma, let

$$N_n := |\{T \in \mathcal{T}'_{(0,0)} : n_e(T) = n\}| \quad (n \geq 0) \quad (5.14)$$

denote the number of non-equivalent contours in  $\mathcal{T}'_{(0,0)}$  that have  $n$  edges of each charge. In Subsection 5.4 we will prove the following exponential bound on  $N_n$ .

**Lemma 40 (Exponential bound)** *Let  $M := |\Delta|$  with  $\Delta$  defined in (5.8) and let  $\tau := \lceil \frac{1}{2}\sigma \rceil$  denote  $\frac{1}{2}\sigma$  rounded up to the next integer. Then*

$$N_n \leq n^{\tau-1}(\tau+1)^{2\tau n} M^{\sigma n} \quad (n \geq 1). \quad (5.15)$$

For our next lemma, we fix a polar function  $L$  satisfying the assumptions of Theorem 9 and we define

$$R := \sum_{s=1}^{\sigma} R_s \quad \text{with} \quad R_s := - \inf_{i \in \Delta} L_s(i) \quad (1 \leq s \leq \sigma), \quad (5.16)$$

and we recall that  $\varepsilon$  and  $\varepsilon_s$  are defined in (5.3). We will prove the following lemma in Subsection 5.3.

**Lemma 41 (Upper bound on the number of edges)** *Each Toom contour  $T \in \mathcal{T}'_{(0,0)}$  satisfies  $n_e(T) \leq (1 + R/\varepsilon)(n_*(T) - 1)$ .*

**Proof of Theorem 9** We use (5.6) which follows from Theorem 27. To prove the stability of  $\Phi^0$ , it is enough to prove that the right-hand-side of (5.6) goes to 0 as  $p \rightarrow 0$ , while by the remarks below (5.6) it suffices to sum over all  $T \in \mathcal{T}'_{(0,0)}$ . By condition (i) of Definition 18 of an embedding, sinks of a Toom contour do not overlap. By condition (i) of Definition 22 of what it means for a Toom contour to be present, each sink corresponds to a space-time point  $(i, t)$  that is defective, meaning that  $\Phi_{i,t}^p = \varphi^0$ , which happens with probability  $p$ , independently for all space-time points. As a result, the probability that a given contour  $T$  is present in  $(\Lambda, \mathcal{H})$  can be estimated from above by  $p^{n_*(T)}$ . By Lemma 41, it follows that

$$\begin{aligned} 1 - \bar{\rho}(p) &\leq \sum_{T \in \mathcal{T}'_{(0,0)}} \mathbb{P}[T \text{ is present in } (\Lambda, \mathcal{H})] \leq \sum_{T \in \mathcal{T}'_{(0,0)}} p^{n_*(T)} = p \sum_{T \in \mathcal{T}'_{(0,0)}} p^{n_*(T)-1} \\ &\leq p \sum_{T \in \mathcal{T}'_{(0,0)}} p^{n_e(T)/(1+R/\varepsilon)} = p \sum_{n=0}^{\infty} N_n p^{n/(1+R/\varepsilon)}, \end{aligned} \quad (5.17)$$

Combining (5.17) and Lemma 40, we see that this sum is finite for  $p$  sufficiently small and hence (by dominated convergence) tends to zero as  $p \rightarrow 0$ . This proves that  $\bar{\rho}(p) \rightarrow 1$  as  $p \rightarrow 0$ . ■

### 5.3 Bounding the edges in terms of the sinks

In this subsection, we prove Lemma 41, which says that for Toom contours in  $\mathcal{T}'_{(0,0)}$ , the number of edges can be bounded in terms of the number of sinks. Before we give the formal proof, we explain the main idea, which is really the central idea behind the proof of Theorem 9 and the definition of Toom contours.

As explained in the previous subsection, the probability that a contour  $T$  is present can be estimated from above by  $p^{n_*(T)}$ , where  $n_*(T)$  is the number of sinks of the Toom contour. Therefore, we can estimate the expected number of Toom contours that is present in  $(\Lambda, \mathcal{H})$  from above by  $\sum_n M_n p^n$ , where  $M_n$  denotes the number of non-equivalent contours in  $\mathcal{T}'_{(0,0)}$  with  $n$  sinks. In general, it is difficult to control the number of contours with a given number of sinks. As shown in Lemma 40, however, we have good control over the number of contours with a given number of edges. Therefore, as we have seen in the proof of Theorem 9, to show that the Peierls sum in (5.6) is small if  $p$  is small, it suffices to have a result like Lemma 41 that bounds the number of edges from above in terms of the number of sinks.

It is precisely here that condition (1.21) of Theorem 9 on the worst-case edge speeds is used. Recall that  $\vec{E}_s$  are the directed edges of charge  $s$ , which are distinguished as in (2.15) into those that come out of a source other than the root (the set  $\vec{E}_s^\circ$ ) and the others (the set  $\vec{E}_s^\bullet$ ). Condition (ii) of Definition 22 says that edges in  $\vec{E}_s^\bullet$  must be embedded at edges of the same charge of the typed dependence graph  $(\Lambda, \mathcal{H})$ . Here  $\vec{H}_s$  is defined in (2.21) where the sets  $A_{s,i}$  are chosen in relation to the polar function  $L$  as in (5.4) and (5.5). The result of all this is that:

- The function  $L_s$  must increase by at least  $\varepsilon_s$  along each edge of charge  $s$ , except for edges that come out of sources other than the root.

Using this, condition (1.21), Lemma 39, and the fact that one edge of each charge originates at each source and one edge of each charge arrives at each sink, we can bound the number of edges in  $\vec{E}_s^\bullet$  in terms of the number of edges in  $\vec{E}_s^\circ$ . Since there are  $|\vec{E}_s^\circ| + 1$  sources and an equal number of sinks, this allows us to bound the total number of edges in terms of the number of sinks.

We now make these ideas precise and prove Lemma 41. We start with a general observation. On any set  $\Lambda$ , we define a *polar function* of *dimension*  $\sigma \geq 2$  to be a function  $L : \Lambda \rightarrow \mathbb{R}^\sigma$  such that

$$\sum_{s=1}^{\sigma} L_s(i) = 0 \quad (i \in \Lambda). \quad (5.18)$$

The following lemma makes a connection between Toom contours and polar functions.

**Lemma 42 (Zero sum property)** *Let  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  be a Toom contour with  $\sigma \geq 2$  charges and let  $L : \Lambda \rightarrow \mathbb{R}^\sigma$  be a polar function of dimension  $\sigma$ . Then*

$$\sum_{s=1}^{\sigma} \sum_{(v,w) \in \vec{E}_s} (L_s(\psi(w)) - L_s(\psi(v))) = 0. \quad (5.19)$$

**Proof** We can rewrite the sum in (5.19) as

$$\sum_{v \in V} \left\{ \sum_{s=1}^{\sigma} \sum_{(v,w) \in \vec{E}_{s,\text{out}}(v)} L_s(\psi(w)) - \sum_{s=1}^{\sigma} \sum_{(u,v) \in \vec{E}_{s,\text{in}}(v)} L_s(\psi(u)) \right\}. \quad (5.20)$$

At internal vertices, the term inside the brackets is zero because the number of incoming edges of each charge equals the number of outgoing edges of that charge. At the sources and sinks,

the term inside the brackets is zero by the defining property (5.18) of a polar function, since there is precisely one outgoing (resp. incoming) edge of each charge.  $\blacksquare$

**Proof of Lemma 41** We trivially “lift” the linear polar function  $L$ , which is defined on  $\mathbb{Z}^d$ , to the space-time set  $\Lambda = \mathbb{Z}^{d+1}$  by setting

$$L_s(i, t) := L_s(i) \quad (i \in \mathbb{Z}^d, t \in \mathbb{Z}). \quad (5.21)$$

Let  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi) = T \in \mathcal{T}'_{(0,0)}$ . We claim that

$$\begin{aligned} L_s(\psi(w)) - L_s(\psi(v)) &\geq \varepsilon_s && \text{if } (v, w) \in \vec{E}_s^\bullet, \\ L_s(\psi(w)) - L_s(\psi(v)) &\geq -R_s && \text{if } (v, w) \in \vec{E}_s^\circ. \end{aligned} \quad (5.22)$$

Indeed, by condition (ii)’ in the definition of the set  $\mathcal{T}'_{(0,0)}$  in (5.7),  $(v, w) \in \vec{E}_s^\bullet$  implies  $\psi(w) = \psi(v) + (j, -1)$  for some  $j \in \Delta_s = \bigcup_{k=1}^m A_{s,k}$ . The linearity of  $L_s$  implies that  $L_s(\psi(w)) - L_s(\psi(v)) = L_s(j)$ , which is  $\geq \varepsilon_s$  for all  $j \in \Delta_s$  by (5.3) and (5.4). The second inequality in (5.22) follows in the same way from condition (iii)’ in (5.7) and (5.16).

By their definition in (2.15) and Lemma 39, we have

$$|\vec{E}_s^\circ| = n_*(T) - 1 \quad \text{and} \quad |\vec{E}_s^\bullet| = n_e(T) - n_*(T) + 1 \quad (1 \leq s \leq \sigma). \quad (5.23)$$

Lemma 42, (5.22), and (5.23) now imply that

$$\begin{aligned} 0 &= \sum_{s=1}^{\sigma} \left( \sum_{(v,w) \in \vec{E}_s^\bullet} (L_s(\psi(w)) - L_s(\psi(v))) + \sum_{(v,w) \in \vec{E}_s^\circ} (L_s(\psi(w)) - L_s(\psi(v))) \right) \\ &\geq \sum_{s=1}^{\sigma} [(n_e(T) - n_*(T) + 1)\varepsilon_s - (n_*(T) - 1)R_s] = \varepsilon n_e(T) - (\varepsilon + R)(n_*(T) - 1), \end{aligned} \quad (5.24)$$

which implies  $n_e(T) \leq (1 + R/\varepsilon)(n_*(T) - 1)$ .  $\blacksquare$

## 5.4 Exponential bounds on the number of contours

In this subsection, we provide the only missing ingredient in the proof of Theorem 9, which is the proof of Lemma 40. If we would be satisfied with just any exponential bound, then the proof could be quite short, but with a view towards Proposition 11 we will argue a bit more carefully to get a sharper bound.

**Proof of Lemma 40** We first consider the case that the number of charges  $\sigma$  is even. Let  $T = (v_\circ, \mathcal{V}, \mathcal{E}, \psi) \in \mathcal{T}'_{(0,0)}$ . Recall that  $(\mathcal{V}, \mathcal{E})$  is a typed directed graph with  $\sigma$  types of edges, that are called charges. In  $(\mathcal{V}, \mathcal{E})$ , all edges point in the direction from the sources to the sinks. We modify  $(\mathcal{V}, \mathcal{E})$  by reversing the direction of edges of the charges  $\frac{1}{2}\sigma + 1, \dots, \sigma$ . Let  $(\mathcal{V}, \mathcal{E}')$  denote the modified graph. In  $(\mathcal{V}, \mathcal{E}')$ , the number of incoming edges at each vertex equals the number of outgoing edges. Since moreover the undirected graph  $(V, E)$  is connected, it is not hard to see<sup>8</sup> that it is possible to walk through the directed graph  $(\mathcal{V}, \mathcal{E}')$  starting from the root using an edge of charge 1, in such a way that each directed edge of  $\mathcal{E}'$  is traversed exactly once.

Let  $l := \sigma n_e(T)$  denote the total number of edges of  $(\mathcal{V}, \mathcal{E}')$  and for  $0 < k \leq l$ , let  $(v_{k-1}, v_k) \in \vec{E}'_{s_k}$  denote the  $k$ -th step of the walk, which has charge  $s_k$ . Write  $\psi(v_k) =: \psi(v_{k-1}) + (\delta_k, \pm 1)$  where  $\delta_k$  is the spatial increment of the  $k$ -th step and  $\pm 1$  is the temporal increment, which is determined by the charge  $s_k$  of the  $k$ -th step: it is -1 for charges  $1, \dots, \frac{1}{2}\sigma$

<sup>8</sup>This is a simple variation of the “Bridges of Königsberg” problem that was solved by Euler.

and 1 otherwise. Let  $k_0, \dots, k_{\sigma/2}$  denote the times when the walk visits the root  $v_\circ$ . We claim that in order to specify  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  uniquely up to equivalence, in the sense defined in Definition 20, it suffices to know the sequences

$$(s_1, \dots, s_l), \quad (\delta_1, \dots, \delta_l), \quad \text{and} \quad (k_0, \dots, k_{\sigma/2}). \quad (5.25)$$

Indeed, the sinks and sources correspond to changes in the temporal direction of the walk which can be read off from the charges. Although the images under  $\psi$  of sources may overlap, we can identify which edges connect to the root, and we also know the increment of  $\psi(v_k) - \psi(v_{k-1})$  in each step, hence all objects in (2.11) can be identified.

The first charge  $s_1$  is 1 and after that, in each step, we have the choice to either continue with the same charge or choose one of the other  $\frac{1}{2}\sigma$  available charges. This means that there are no more than  $(\frac{1}{2}\sigma + 1)^{l-1}$  possible ways to specify the charges  $(s_1, \dots, s_l)$ . Recalling  $M = |\Delta| = |\bigcup_{s=1}^\sigma \bigcup_{k=1}^m A_{s,k}|$ , we see that there are no more than  $M^l$  possible ways to specify the spatial increments  $(\delta_1, \dots, \delta_l)$ . Since  $k_0 = 0, k_{\sigma/2} = l$ , we can roughly estimate the number of ways to specify the visits to the root from above by  $n^{\sigma/2-1}$ . Recalling that  $l = \sigma n_e(T)$ , this yields the bound

$$N_n \leq n^{\sigma/2-1} (\frac{1}{2}\sigma + 1)^{\sigma n-1} M^{\sigma n}. \quad (5.26)$$

This completes the proof when  $\sigma$  is even.

When  $\sigma$  is odd, we modify  $(\mathcal{V}, \mathcal{E})$  by doubling all edges of charge  $\sigma$ , i.e., we define  $(\mathcal{V}, \mathcal{F})$  with

$$\vec{F} = (\vec{F}_1, \dots, \vec{F}_{\sigma+1}) := (\vec{E}_1, \dots, \vec{E}_\sigma, \vec{E}_\sigma), \quad (5.27)$$

and next we modify  $(\mathcal{V}, \mathcal{F})$  by reversing the direction of all edges of the charges  $[\frac{1}{2}\sigma] + 1, \dots, \sigma + 1$ . We can define a walk in the resulting graph  $(\mathcal{V}, \mathcal{F}')$  as before and record the charges and spatial increments for each step, as well as the visits to the root. In fact, in order to specify  $(v_\circ, \mathcal{V}, \mathcal{E}, \psi)$  uniquely up to equivalence, we do not have to distinguish the charges  $\sigma$  and  $\sigma + 1$ . Recall that edges of the charges  $\sigma$  and  $\sigma + 1$  result from doubling the edges of charge  $\sigma$  and hence always come in pairs, connecting the same vertices. Since sinks do not overlap and internal vertices of a given charge do not overlap, and since we traverse edges of the charges  $\sigma$  and  $\sigma + 1$  in the direction from the sinks towards the sources, whenever we are about to traverse an edge that belongs to a pair of edges of the charges  $\sigma$  and  $\sigma + 1$ , we know whether we have already traversed the other edge of the pair. In view of this, for each pair, we only have to specify the spatial displacement at the first time that we traverse an edge of the pair. Using these considerations, we arrive at the bound

$$N_n \leq n^{[\sigma/2]-1} ([\frac{1}{2}\sigma] + 1)^{(\sigma+1)n-1} M^{\sigma n}. \quad (5.28)$$

■

## 5.5 Some bounds for Toom cycles

With Theorem 9 proved, we start to prepare for the proof of Proposition 11. In the present subsection, we prove more precise versions of Lemmas 40 and 41 that hold only for Toom cycles and that will help us to get a better bound for the critical noise parameter  $p_c$  of the cellular automaton defined by the map  $\varphi^{\text{coop}}$ . Recall that at the end of Subsection 5.1 we denoted the set of non-isomorphic Toom cycles rooted at  $(0, 0)$  by  $\bar{\mathcal{T}}_{(0,0)}$  and we wrote  $\bar{\mathcal{T}}'_{(0,0)}$  for the set of  $T \in \bar{\mathcal{T}}_{(0,0)}$  that satisfy (5.9).

Similarly to (5.14), we let

$$\bar{N}_n := |\{T \in \bar{\mathcal{T}}'_{(0,0)} : n_e(T) = n\}| \quad (n \geq 0) \quad (5.29)$$

denote the number of non-isomorphic Toom cycles in  $\bar{\mathcal{T}}'_{(0,0)}$  that have  $n$  edges of each charge. We then have the following analogue of Lemma 40.

**Lemma 43 (Exponential bound for  $\sigma = 2$ )** *Let  $M_s := |\Delta_s|$  ( $s = 1, 2$ ) with  $\Delta_s$  defined in (5.8). Then*

$$\bar{N}_n \leq \frac{1}{2}(4M_1M_2)^n \quad (n \geq 1). \quad (5.30)$$

**Proof** The proof goes along the same lines as that of Lemma 40 for the case  $\sigma$  is even. Observe that for  $\sigma = 2$ , the walk visits the root 0 twice:  $k_0 = 0, k_1 = l$ , where  $l$  is the total number of edges of the cycle. Thus  $(k_0, k_1)$  is deterministic, and we only need to specify the sequences

$$(s_1, \dots, s_l), \quad (\delta_1, \dots, \delta_l). \quad (5.31)$$

Note that in this case, these sequences determine the Toom cycle up to isomorphism and not only up to equivalence as in the proof of Lemma 40. The first charge  $s_1$  is 1 and after that, in each step, we have the choice to either continue with the same charge or choose charge 2. This means that there are no more than  $2^{l-1}$  possible ways to specify the charges  $(s_1, \dots, s_l)$ . Once we have done that, by condition (iii)" of (5.9), we know for each  $0 < k \leq l$  whether the spatial increment  $\delta_k$  is in  $\Delta_1$  or  $\Delta_2$ . Recalling  $M_s = |\Delta_s|$  ( $s = 1, 2$ ) and using the fact that  $|\vec{E}_1| = |\vec{E}_2| = n_e(T) = l/2$ , we see that there are no more than  $M_1^{l/2} \cdot M_2^{l/2}$  possible ways to specify  $(\delta_1, \dots, \delta_l)$ . This yields the bound

$$\bar{N}_n \leq 2^{2n-1} M_1^n \cdot M_2^n. \quad (5.32)$$

■

From now on, we fix a polar function  $L$  of dimension 2 satisfying the assumptions of Theorem 9. In analogy with (5.16), but with a view towards (5.9) which in the present context replaces (5.7), we define

$$\bar{R} := \sum_{s=1}^2 \bar{R}_s \quad \text{with} \quad \bar{R}_1 := - \inf_{i \in \Delta_2} L_1(i) \quad \text{and} \quad \bar{R}_2 := - \inf_{i \in \Delta_1} L_2(i). \quad (5.33)$$

The following lemma is similar to Lemma 41.

**Lemma 44 (Upper bound on the number of edges for  $\sigma = 2$ )** *Let  $\varepsilon$  be defined in (1.21) and let  $\bar{R}$  be defined in (5.33). Then each  $T \in \bar{\mathcal{T}}'_{(0,0)}$  satisfies  $n_e(T) \leq (1 + \bar{R}/\varepsilon)(n_*(T) - 1)$ .*

**Proof** The proof is the same as that of Lemma 41, with the only difference that condition (iii)" of (5.9) allows us to use  $\bar{R}_s$  instead of  $R_s$  ( $s = 1, 2$ ) as upper bounds. ■

## 5.6 Finiteness of the Peierls sum

We continue our preparations for the proof of Proposition 11. Our aim is to derive a lower bound  $p_*$  on the critical noise parameter  $p_c$  defined in (1.7), which requires us to prove that  $\bar{\rho}(p) > 0$  for all  $p < p_*$ . By (5.6), we have  $\bar{\rho}(p) > 0$  as soon as the Peierls sum on the right-hand side of (5.6) is less than one. In the present subsection, we will prove that in fact it (more or less) suffices to show that the Peierls sum is finite. This will not only lead to slightly better bounds but also simplify our calculations later. Similar results, which say that finiteness of the Peierls sum already implies a phase transition, have been proved before. For percolation, the argument is quite simple [Dur88, Section 6a] but for other models such results can be a bit harder to obtain [KSS14].

We will work in the set-up of Subsection 5.1, but specialised to the case  $m = 1$ , which means that  $\Phi^0$  is a deterministic monotone cellular automaton. To simplify notation, we set  $\varphi := \varphi_1$  and  $A_s := A_{s,1}$ . We recall from (5.4) that  $A_s$  is chosen such that

$$\sup_{A \in \mathcal{A}(\varphi)} \inf_{i \in A} L_s(i) = \inf_{i \in A_s} L_s(i) \quad (1 \leq s \leq \sigma), \quad (5.34)$$

and that by (1.15) and (5.3)

$$\varepsilon = \sum_{s=1}^{\sigma} \varepsilon_s \quad \text{with} \quad \varepsilon_s = \varepsilon_{\varphi}(L_s) = \inf_{i \in A_s} L_s(i) \quad (1 \leq s \leq \sigma). \quad (5.35)$$

In our present setting, the definition of the sets  $\Delta_s$  and  $\Delta$  from (5.8) simplifies to

$$\Delta_s = A_s \quad (1 \leq s \leq \sigma) \quad \text{and} \quad \Delta = \bigcup_{s=1}^{\sigma} A_s. \quad (5.36)$$

As in Subsection 5.1, we let  $\mathcal{T}'_{(0,0)}$  denote the set of Toom contours rooted at  $(0,0)$  that satisfy (5.7). In the special case  $\sigma = 2$ , we let  $\bar{\mathcal{T}}'_{(0,0)}$  denote the set of Toom cycles rooted at  $(0,0)$  that satisfy (5.9). As in (5.11) we let  $n_*(T)$  denote the number of sinks of a Toom contour  $T$ , which equals the number of sources. Here is the main result of this subsection.

**Proposition 45 (Finiteness of the Peierls sum)** *Assume that  $\varepsilon > 0$ . Then the condition*

$$\sum_{T \in \mathcal{T}'_{(0,0)}} p^{n_*(T)} < \infty \quad (5.37)$$

*implies that  $\bar{\rho}(p) > 0$ . If  $\sigma = 2$ , then the same conclusion can be drawn if in (5.37) the sum over  $\mathcal{T}'_{(0,0)}$  is replaced by the sum over  $\bar{\mathcal{T}}'_{(0,0)}$ .*

Proposition 45 actually stops short of what we promised, since the sum in (5.37) is only an upper bound for the Peierls sum on the right-hand side of (5.6). For our purposes, the statement of Proposition 45 will be sufficient, however.

The proof of Proposition 45 needs some preparations. Recall that  $\Lambda = \mathbb{Z}^{d+1}$  and that the function  $\kappa$  is defined in (5.1). In the present setting,  $(\kappa(i,t))_{(i,t) \in \Lambda}$  are i.i.d.  $\{0,1\}$ -valued random variables with  $\mathbb{P}[\kappa(i,t) = 1] = p$ , and

$$\Lambda_0 := \{(i,t) \in \Lambda : \kappa(i,t) = 0\} \quad \text{and} \quad \Lambda_{\bullet} := \{(i,t) \in \Lambda : \kappa(i,t) = 1\}. \quad (5.38)$$

Recall that sites in  $\Lambda_0$  are called *defective*. In our present setting, the definition of the typed dependence graph  $(\Lambda, \mathcal{H})$  simplifies to

$$\vec{H}_s := \{((i,t), (i+j, t-1)) : (i,t) \in \Lambda_{\bullet}, j \in A_s\} \quad (1 \leq s \leq \sigma). \quad (5.39)$$

It will be convenient to define a modified typed dependence graph  $(\Lambda, \mathcal{H}^<)$  that has no defective sites  $(i,t)$  with time coordinates  $t > 0$ . Formally, we define  $\Lambda_0^< := \{(i,t) \in \Lambda_0 : t \leq 0\}$ ,  $\Lambda_{\bullet}^< := \Lambda \setminus \Lambda_0^<$ , and we define  $\mathcal{H}^< = (\vec{H}_1^<, \dots, \vec{H}_{\sigma}^<)$  as in (5.39) but with  $\Lambda_{\bullet}$  replaced by  $\Lambda_{\bullet}^<$ . We let  $\Psi^<$  denote the monotone cellular automaton associated with  $(\Lambda, \mathcal{H}^<)$ , in the sense of Definition 21, and we let  $\bar{Y}^<$  denote the maximal trajectory of  $\Psi^<$ .

**Lemma 46 (Presence of a large contour)** *Fix  $j_s \in A_s$  ( $1 \leq s \leq \sigma$ ) and  $r \in \mathbb{N}$ . Let  $C_r \subset \mathbb{Z}^d$  with  $r \in \mathbb{N}$  be inductively defined by*

$$C_0 := \{0\} \quad \text{and} \quad C_{r+1} := \{i + j_s : i \in C_r, 1 \leq s \leq \sigma\} \quad (r \geq 0). \quad (5.40)$$

*Then on the event that  $\bar{Y}^<(i,0) = 0$  for all  $i \in C_r$ , there is a Toom contour  $(v_{\circ}, \mathcal{V}, \mathcal{E}, \psi)$  rooted at  $(0,r)$  present in  $(\Lambda, \mathcal{H}^<)$ . If  $\sigma = 2$ , then a Toom cycle rooted at  $(0,r)$  is present in  $(\Lambda, \mathcal{H}^<)$ .*



**Proof** Since there are no defective sites with positive time coordinates, by Definition 21 and (5.39) we have

$$\bar{Y}^{\langle}(i, t) = \bigvee_{s=1}^{\sigma} \bigwedge_{j \in A_s} Y^{\langle}(i + j, t - 1) \quad (i \in \mathbb{Z}^d, t > 0). \quad (5.41)$$

Using this and the assumption that  $\bar{Y}^{\langle}(i, 0) = 0$  for all  $i \in C_r$  we see by induction that  $\bar{Y}^{\langle}(i, t) = 0$  for all  $i \in C_{r-t}$  ( $0 \leq t \leq r$ ) and hence in particular  $\bar{Y}^{\langle}(0, r) = 0$ . The claim now follows from Theorems 23 and 26.  $\blacksquare$

**Lemma 47 (Many sinks)** *Let  $\varepsilon$  and  $R$  be defined as in (1.21) and (5.16). Assume that  $T$  is a Toom contour rooted at  $(0, r)$  that is present in  $(\Lambda, \mathcal{H}^{\langle})$ . Then  $n_*(T) \geq r\varepsilon/R + 1$ .*

**Proof** This follows from an argument similar to the proof of Lemma 41. Since  $(\Lambda, \mathcal{H}^{\langle})$  has no defective sites with positive time coordinates, any Toom contour  $T = (v_{\circ}, \mathcal{V}, \mathcal{E}, \psi)$  that is rooted at  $(0, r)$  and present in  $(\Lambda, \mathcal{H}^{\langle})$  must satisfy  $|\bar{E}_s^{\bullet}| \geq r$  ( $1 \leq s \leq \sigma$ ), so a calculation as in (5.24) gives

$$0 \geq \sum_{s=1}^{\sigma} [r\varepsilon_s - (n_*(T) - 1)R_s] = r\varepsilon - (n_*(T) - 1)R. \quad (5.42)$$

**Proof of Proposition 45** Let  $\Psi$  denote the monotone cellular automaton associated with  $(\Lambda, \mathcal{H})$ . Let  $\bar{X}^p$  denote the maximal trajectory of  $\Phi^p$  and let  $\bar{Y}$  denote the maximal trajectory of  $\Psi$ . As in the proof of Theorem 27 we see that  $\bar{Y} \leq \bar{X}^p$  (pointwise). Moreover, since  $(\Lambda, \mathcal{H})$  and  $(\Lambda, \mathcal{H}^{\langle})$  agree up to time zero, it is easy to see that  $\bar{Y}(i, t) = \bar{Y}^{\langle}(i, t)$  for all  $t \leq 0$ . Let  $M_n$  denote the number of non-equivalent contours in  $\mathcal{T}'_{(0,0)}$  with  $n$  sinks. Lemmas 46 and 47 tell us that

$$\mathbb{P}[\bar{Y}^{\langle}(i, 0) = 0 \ \forall i \in C_r] \leq \sum_{n=r\varepsilon/R+1}^{\infty} M_n p^n, \quad (5.43)$$

and the same is true with  $\bar{Y}^{\langle}$  replaced by  $\bar{Y}$ , since they are equal at time zero. Using (5.37) and the assumption that  $\varepsilon > 0$ , we see that we can choose  $r$  large enough such that the probability in (5.43) is less than one. It follows that

$$0 < \mathbb{P}[\bar{Y}(i, 0) = 1 \text{ for some } i \in C_r] \leq |C_r| \cdot \mathbb{P}[\bar{Y}(0, 0) = 1] \leq |C_r| \cdot \mathbb{P}[\bar{X}^p(0, 0) = 1], \quad (5.44)$$

proving that  $\bar{\rho}(p) > 0$ . The proof for Toom cycles is the same.  $\blacksquare$

## 5.7 Bounds on the critical noise parameter

In this subsection we prove Proposition 11. We work in the set-up of Subsection 5.1, specialised to the case  $m = 1$ , as summarised in Subsection 5.6.

**Proof of Proposition 11** First we consider the cellular automaton that applies the map  $\varphi^{\text{coop}}$  (recall (1.11)) at each space-time point. We chose  $\sigma := 2$ , and the polar function

$$L_1(z) := -z_1 - z_2, \quad L_2(z) := z_1 + z_2 \quad (z_1, z_2) \in \mathbb{R}^2. \quad (5.45)$$

Recalling the minimal one-sets of  $\varphi^{\text{coop}}$  from (1.12), we then choose  $A_1, A_2 \in \mathcal{A}(\varphi^{\text{coop}})$  satisfying (5.34) by setting  $A_1 := \{(0, 0)\}$ , and  $A_2 := \{(0, 1), (1, 0)\}$ . This has the result that the constants from (1.21), (5.16) and (5.33) are given by  $\varepsilon = 1$ ,  $R = 1$  and  $\bar{R} = 1$ . We first give a

bound using Toom contours. Applying Lemma 40 with  $M = 3$ ,  $\sigma = 2$  and  $\tau = 1$ , the Peierls bound (5.17) gives

$$1 - \bar{\rho}(p) \leq \sum_{T \in \mathcal{T}'_{(0,0)}} p^{n_*(T)} \leq p \sum_{n=0}^{\infty} N_n p^{n/(1+R/\varepsilon)} \leq p \sum_{n=0}^{\infty} 2^{2\tau n} 3^{2n} p^{n/2}. \quad (5.46)$$

By Proposition 45, to prove that  $\bar{\rho}(p) > 0$ , it actually suffices to prove that the right-hand side of (5.46) is finite, which happens when  $36p^{1/2} < 1$ , leading to the bound  $p_c \geq 36^{-2}$ .

Since  $\sigma = 2$ , we can use Toom cycles instead. Using (5.6), Lemma 43 with  $M_1 = 1$ ,  $M_2 = 2$ , and Lemma 44, we find that

$$1 - \bar{\rho}(p) \leq \sum_{T \in \mathcal{T}'_{(0,0)}} p^{n_*(T)} \leq p \sum_{n=0}^{\infty} \bar{N}_n p^{n/(1+\bar{R}/\varepsilon)} \leq p + \frac{1}{2}p \sum_{n=1}^{\infty} 8^n p^{n/2}. \quad (5.47)$$

Again, by Proposition 45, it suffices to prove that the right-hand side is finite, which happens when  $8p^{1/2} < 1$ , leading to the bound  $p_c \geq 1/64$ .

Now consider the cellular automaton that applies the map  $\varphi^{\text{NEC}}$  (recall (1.11)) at each space-time point. We chose  $\sigma := 3$ , and the polar function

$$L_1(z_1, z_2) := -z_1, \quad L_2(z_1, z_2) := -z_2, \quad L_3(z_1, z_2) := z_1 + z_2. \quad (5.48)$$

Recalling the minimal one-sets of  $\varphi^{\text{NEC}}$  from (1.12) we choose  $A_1, A_2, A_3 \in \mathcal{A}(\varphi^{\text{NEC}})$  satisfying (5.34) by setting  $A_1 := \{(0, 0), (0, 1)\}$ ,  $A_2 := \{(0, 0), (1, 0)\}$ , and  $A_3 := \{(0, 1), (1, 0)\}$ . This has the result that the constants from (1.21) and (5.16) are given by  $\varepsilon = 1$  and  $R = 2$ . Applying Lemma 40 with  $M = 3$ ,  $\sigma = 3$  and  $\tau = 2$ , the Peierls bound (5.17) gives

$$1 - \bar{\rho}(p) \leq \sum_{T \in \mathcal{T}'_{(0,0)}} p^{n_*(T)} \leq p \sum_{n=0}^{\infty} n 3^{4n} 3^{3n} p^{n/3}. \quad (5.49)$$

By Proposition 45, it suffices to prove that the right-hand side is finite, which happens when  $3^7 p^{1/3} < 1$ , leading to the bound  $p_c \geq 3^{-21}$ .  $\blacksquare$

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